Lecture 10

Today:	IIR filter implementation (ctd) DFT and FFT Spectral leakage FFT scaling
Announcements:	Fri Oct 3 office hours cancelled Extra office hours today 12-1:30PM Deadline for parts orders is Fri Oct 10 Hwk 5 due on thurs Oct 16 Midterm exam on thurs Oct 23. Coverage: hwks 1-5, labs 1-6, lectures 1-12.

Please keep the lab clean and organized.

Last one out should close the lab door!!!!

In mathematics you don't understand things, you just get used to them. — John von Neumann

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IIR Canonical direct form 2



a) Non-canonical Direct Form 2.

b) DF2 in canonical form.

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Implementing a biquad cascade IIR filter

The implementation steps are:

- ▶ Factor the transfer function into pole and zero pairs.
- ▶ Choose a biquad architecture, e.g., DF2, TDF2.
- ▶ Relate the biquad coefficients to the chosen architecture coefficients.
- ► Order the poles and the zeros to control internal resonance levels.
- Distribute the gain between the biquad sections.
- Normalize biquad coefficients if necessary
- ▶ Program and get to work.
- ► Test.

Pole - zero pairing example



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Pole - zero pairing example



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DF2 biquad cascade



The above block diagram shows a cascade of four DF2 second order biquad sections. This can be used to implement an eighth order lowpass filter.

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TDF2 biquad cascade



The above block diagram shows a cascade of four TDF2 second order biquad sections. This can be used to implement an eighth order lowpass filter.

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The DF2 and TDF2 biquad sections



(a) Direct form type 2 biquad section. (b) Transposed direct form 2 biquad section.

$$H(z) = \frac{b_0 + b_1 z^{-1} + b_2 z^{-2}}{1 + a_1 z^{-1} + a_2 z^{-2}}$$

The most commonly used biquad is direct form 2. We need to

analyze the transfer function magnitudes between input and internal states in addition to between input and output.

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DF2 biquad cascade transfer functions



Write $H_i = \frac{Y_i}{X_i}$, $H_{i1} = \frac{W_{i1}}{X_i}$ and $H_{i2} = \frac{W_{i2}}{X_i}$. Because of our choice of the L_{∞} norm we are interested in the magnitudes of the input to delay stage filter functions:

section 1	H_{11}	H_1
section 2	$H_1 H_{21}$	H_1H_2
section 3	$H_1 H_2 H_{31}$	$H_1H_2H_3$
section 4	$H_1 H_2 H_3 H_{41}$	$H_1H_2H_3H_4$

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TDF2 biquad cascade transfer functions



Write $H_i = \frac{Y_i}{X_i}$, $H_{i1} = \frac{W_{i1}}{X_i}$ and $H_{i2} = \frac{W_{i2}}{X_i}$. Because of our choice of the L_{∞} norm we are interested in the magnitudes of the input to delay stage filter functions:

section 1	H_{11}	H_{12}	H_1
section 2	$H_1 H_{21}$	$H_1 H_{22}$	H_1H_2
section 3	$H_1 H_2 H_{31}$	$H_1 H_2 H_{32}$	$H_1H_2H_3$
section 4	$H_1 H_2 H_3 H_{41}$	$H_1 H_2 H_3 H_{42}$	$H_1H_2H_3H_4$

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Filter input-to-delay stage TFs: DF2 - max gain



By nominally scaling the input by 4 we can avoid overflow in this realization. If a 12-bit converter is being used and a 16-bit word size, this is no great loss.

NB: this transfer function isn't the one used in lab.

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Biquad input-to-delay stage TFs: TDF2



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Biquad input-to-delay stage TFs: TDF2 (ctd)



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Filter input-to-delay stage TFs: TDF2



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Filter input-to-delay stage TFs: TDF2 (ctd)



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Filter input-to-delay stage TFs: TDF2 - max gain



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Overflow issues for biquad coefficients

Consider the biquad

$$H(z) = \frac{b_0 + b_1 z^{-1} + b_2 z^{-2}}{1 + a_1 z^{-1} + a_2 z^{-2}} = \frac{b_0 z^2 + b_1 z + b_2}{z^2 + a_1 z + a_2}.$$

The poles of H(z) are determined by $z^2 + a_1 z + a_2$. Assume a complex valued pole pair, $p_1 = re^{j\theta}$ and $p_2 = re^{-j\theta}$.

$$(z - p_1)(z - p_2) = z^2 - 2r\cos(\theta)z + r^2 = z^2 + a_1z + a_2.$$

In order for the filter to be (conditionally) stable the biquad poles have to be (on or) within the unit circle.

Because $0 \le r \le 1$ we have that $0 \le a_2 \le 1$ and $-2 < a_1 \le 2$. In addition, $a_2 \ge a_1^2/4$.

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Overflow issues for biquad coefficients

We will be using Q15 numeric format values in the C5515 and DE2-70. The magnitude of the a_1 value can be greater than 1 (but less than 2). We need to worry about this.

There also may be scaling concerns with the b coefficient values as well. One needs to stay alert.

Occasionally there are b_i values with magnitude greater than 1. Large b values can be handled by scaling all of the b coefficients. This affects only the gain through the system.

Scaling cannot be applied to the *a* values without changing the shape of the transfer function. (Recall $a_0 = 1$ requirement). Alternatives:

- 1. Implement each multiply $a_i x[n]$ as $((a_i/2)x[n])2$.
- 2. Use Q(14) representation.

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Coefficient scaling possibilities

If we divide the a's by k we need to multiply the sum by k. Use Q14 data and Q14 coefficients? Q14×Q14 gives Q28. To make Q28 into Q14 shift left 2 then truncate.

Use Q15 data and Q14 coefficients? Q15×Q14 gives Q29. To make Q29 into Q15 shift left 2 and truncate.

Use Q15 data and Q15 coefficients? Q15 \times Q15 gives Q30. To make Q30 into Q15 shift left 1 and truncate. Need to remember to round before truncating.



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Do you need to normalize any coefficients?





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Utility of DFT and FFT

- ▶ Spectrum analysis:
 - ▶ Measure frequency response of an analog or digital filter
 - ▶ Measure frequency content of a signal (speech, audio, rf, etc)
 - Estimate magnitude or phase of an unknown channel.
- Spectrum synthesis: vocoder, voice synthesis, voice scrambler, voice coding, audio effects.
- ▶ Implementing filters in frequency domain with DFT or FFT:
 - ► An FIR LTI filter can be implemented by convolving input with impulse response h[n]

$$y[n] = \sum_{k=0}^{N-1} h[k]x[n-k],$$

• ... or by IDFT of the product of $H(k) = DFT_k(h[n])$ and $DFT_k(x[n])$

$$Y[k] = H[k]X[k], \quad k = 0, \dots, N-1$$

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Filter implementation in time or in frequency

$$x[n] \longrightarrow h[n] \longrightarrow y[n]$$

$$H[k] = \text{DFT}_{k}(h[n]) = \sum_{n=0}^{N-1} h[n]e^{-j2\pi \frac{k}{N}n}$$



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Discrete Fourier Transform

Available: N time samples $x[0], \ldots, x[N-1]$. DFT is defined for $k = 0, \ldots, N-1$

$$X_{DFT}(k) = \text{DFT}_k(x[n]) = \sum_{n=0}^{N-1} x[n] e^{-j2\pi \frac{k}{N}n}$$

IDFT recovers time samples from DFT

$$x[n] = \text{IDFT}_{n}(X_{DFT}(k)) = \frac{1}{N} \sum_{n=0}^{N-1} X_{DFT}(k) e^{j2\pi \frac{k}{N}n}$$

Rectangular form of DFT

$$X_{DFT}[k] = \sum_{n=0}^{N-1} x[n] \cos(2\pi nk/N) - j \sum_{n=0}^{N-1} x[n] \sin(2\pi nk/N)$$

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64 point DFT of a discrete time signal



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Some properties of DFT

 $X[k] = \text{DFT}_k(x[n]), \{x[n]\}_{n=0}^{N-1}$ is real valued, N is even integer.

- ▶ Conjugate symmetry: DFT satisfies $X[N-k] = X^*[k]$
- Magnitude symmetry:

$$|X[N/2 + k]| = |X[N/2 - k]|, \quad k = 0, \dots, N/2$$
$$|X[N - k]| = |X[k]|, \quad k = 0, \dots, N/2$$

▶ Phase anti-symmetry:

$$\begin{split} \arg(X[N/2+k]) &= -\arg(X[N/2-k]), \ k = 0, \dots, N/2\\ \arg(X[N-k]) &= -\arg(X[k]), \ k = 0, \dots, N/2\\ \left(\arg X = \texttt{angle}(X) = \texttt{atan}\left(\frac{Im(X)}{Re(X)}\right)\right) \end{split}$$

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The Fast Fourier Transform (FFT)

Q. How many MAC's does the DFT consume?

$$X[k] = \sum_{n=0}^{N-1} x[n] e^{-j2\pi kn/N}, \quad k = 0, 1, \dots, N-1.$$

A. The nominal computational cost is N^2 complex multiply-adds.

Any algorithm that significantly reduces this number can be considered as being *fast*.

There are many *fast* DFT algorithms. Some algorithms are faster than others under different circumstances.

We will only cover the original Cooley-Tukey FFT.

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Roots of unity, powers of $W_N = e^{-j2\pi/N}$.

Symmetry of the sine and cosine.

Index mappings.

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The decimation-in-time radix-2 FFT

- N is assumed to be an integer power of 2.
- Divide $\{x[n]\}_{n=0}^{N-1}$ into $\{x[2n]\}_{n=0}^{N/2-1}$ and $\{x[2n+1]\}_{n=0}^{N/2-1}$
- ▶ Form the DFT of each set and combine results to form N value DFT.
- ▶ Repeat the procedure on each of the N/2-point DFTs.
- ► And so on.

The resulting nominal complex mult-add count is $N \times \log_2(N)$

N	$\log_2(N)$	$N \times \log_2(N)$	N^2
8	3	24	64
16	4	64	256
32	5	160	1024
64	6	384	4096
128	7	896	16384
256	8	2048	65536
512	9	4068	262144
1024	10	10249	1048576

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Easy when you see how it's done

Start with the forward transform equation

$$X[k] = \sum_{n=0}^{N-1} x[n] e^{-j2\pi kn/N}, \quad k = 0, 1, \dots, N-1.$$

Express even n as 2p and odd n as 2q + 1

$$\begin{split} X[k] &= \sum_{p=0}^{N/2-1} x[2p] e^{-j2\pi k 2p/N} + \sum_{q=0}^{N/2-1} x[2q+1] e^{-j2\pi k (2q+1)/N} \\ &= \sum_{p=0}^{N/2-1} x[2p] e^{-j2\pi k p/(N/2)} + e^{-j2\pi k/N} \sum_{q=0}^{N/2-1} x[2q+1] e^{-j2\pi k q/(N/2)} \,. \end{split}$$

Symmetry relation for k = 1, ..., N/2 - 1:

$$X[k+N/2] = \sum_{p=0}^{N/2-1} x[2p]e^{-j2\pi kp/(N/2)} - e^{-j2\pi k/N} \sum_{q=0}^{N/2-1} x[2q+1]e^{-j2\pi kq/(N/2)}$$

Have saved half the computations (N/2). Repeat the process $\log_2 N$ times. EECS 452 - Fall 2014 Lecture 10 - Page 29/46 Tuesday - 10/2/2014

Example: 8-point radix-2 DIT FFT

DFT MACS: $N^2 = 64$ FFT MACS: $N \log(N) = 24$ W_N is defined as $e^{-j2\pi/N}$.

N = 8 for this example.

Arrows indicate multiplication.

Nodes represent summing.



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Example: subdivide again



Multiplication by -1 is the result of a 2-point DFT. Values flow left to right.

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Modified 8-point radix-2 FFT diagram

W is defined as $e^{-j2\pi/N}$. N = 8 for this example. Arrows indicate multiplication. Nodes represent summing. Multiplication by -1 is trivial. Values flow left to right.



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Reordering of DIT FFT input/output values

Can reorder the flow graph so that the input is in normal order and the output is in bit reverse order.



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Using FFT for signal analysis

- ▶ The k-th coefficient of the N-point FFT of x[n] is a sample of the DTFT of x[n] at digital frequency f = k/N.
- If x[n] are time samples $x(nT_s)$ of a continuous time signal x(t) then DTFT is an approximation to the finite time FT of x(t) over the time window $t \in [0, (N-1)T_s)$.
- ▶ There are several issues that need to be addressed
 - Spectral leakage
 - Spectral resolution
 - Time varying spectra
- ▶ To build intuition we start by considering an example: DFT of sinusoidal signal.

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DFT of a sinusoid at frequency $f_c = m/N$

DFT of sinusoid $x[n] = \cos(2\pi f_c n + \phi)$?

Assume sinusoidal frequency satisfies $f_c = m/N$ for integer $m \in \{0, \dots, N/2\}$

Use Euler formula: $\cos(\theta) = (e^{j\theta} + e^{-j\theta})/2$

$$X_{DFT}(k) = \frac{e^{j\phi}}{2} \underbrace{\sum_{n=0}^{N-1} e^{-j2\pi \frac{k-m}{N}n}}_{N\Delta[k-m]} + \frac{e^{-j\phi}}{2} \underbrace{\sum_{n=0}^{N-1} e^{-j2\pi \frac{k+m}{N}n}}_{N\Delta[N-k-m]}$$
$$= \begin{cases} \frac{Ne^{j\phi}}{2}, & k=m\\ \frac{Ne^{-j\phi}}{2}, & k=N-m \end{cases}$$

 $(\Delta[n] \text{ is kronecker delta function})$

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DFT of single sinusoid at arbitrary freq

DFT of sinusoid $x[n] = cos(2\pi f_c n + \phi)$?

Assume sinusoidal frequency **does not** satisfy $f_c = m/N$ for integer $m \in \{0, \dots, N/2\}$

$$X_{DFT}(k) = \frac{e^{j\phi}}{2} \underbrace{\sum_{n=0}^{N-1} e^{-j2\pi \frac{k-Nf_c}{N}n}}_{\neq N\Delta[k-m]} + \frac{e^{-j\phi}}{2} \underbrace{\sum_{n=0}^{N-1} e^{-j2\pi \frac{k+Nf_c}{N}n}}_{\neq N\Delta[N-k-m]}$$

This is the **leakage** phenomenon and it occurs when $f_c \neq m/N$.

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DFT: a sinusoid $f_c = 0.25 = m/N$, m = 4, N = 64



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DFT: single sinusoid $f_c = 4.5/N$, not an integer/N



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DFT: two orthogonal sinusoids $f_{ci} = m_i/N$, $m_1 = 4$, $m_2 = 5$, N = 64



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DFT: two non-orthogonal sinusoids

 $f_{ci} = m_i/N$, $m_1 = 4$, $m_2 = 4.5$, N = 64



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Two orthogonal sinusoids $f_{ci} = m_i/N$, $m_1 = 4$,

 $m_2 = 5$, N = 64



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Two nonorthogonal sinusoids $f_{ci} = m_i/N$,

 $m_1 = 4$, $m_2 = 4.5$, N = 64



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FFT input scaling

Consider using standard 16-bit Q15 number representation in FFT as in AIC3204.

Let input to the FFT be the cosine signal

$$\cos(2\pi f_c n) = \frac{e^{j2\pi f_c n} + e^{-j2\pi f_c n}}{2}, \ n = 0, \dots, N-1$$

Overflow problem 1: The gain at the f_c frequency (assuming it matches some analysis frequency m/N) is N/2. If a 1024 point transform is taken then the result might require 10-1+16 = 25 bits.

Overflow problem 2: A complex input with 16 bit Q15 real and imaginary parts can overflow if a phase rotation occurs. For example, 1 + j1 can rotate to 1.414 + j0 creating an overflow in the Q15 real part.

This is why in lab 6 you will be implementing 32 bit precision FFT's.

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• Normalization

Consider a Q15 sinewave input having amplitude 1. Using 1/N scaling on the forward transform, the magnitude of the FFT output will be capped at 1/2.

• Distribute normalization over each of the $\log_2(N)$ layers of FFT

Assume $N = 2^n$ is a power of two, $n = \log_2 N$ an integer. Then can apply a scale factor of 1/2 to each layer of the FFT. The net effect will be to scale the FFT operation by 1/N.

Summary of what we covered today

- ▶ The DFT and FFT
- Finite precision and scaling issues for FFT
- Spectral leakage

- "Applied signal processing," Dutoit and Marques, 2010.
- "Digital signal processing," Proakis and Manolakis, 3rd Edition.
- "Understanding digital signal processing," R. Lyons, 2004.