

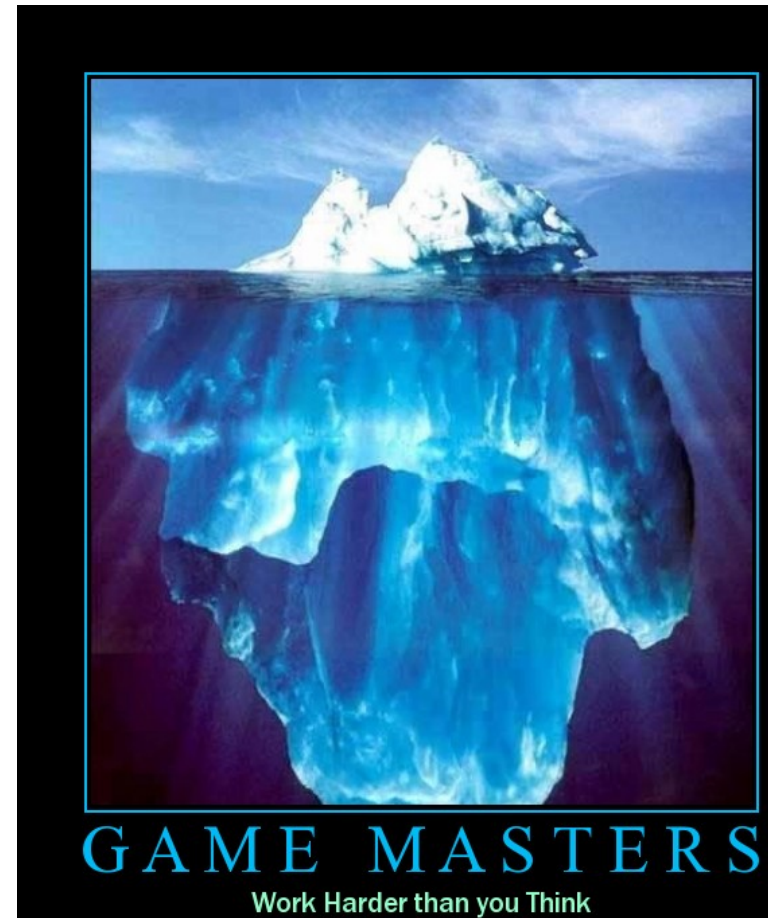


# Functional Programming

## Introduction To Cool

# Cunning Plan

- ML Functional Programming
  - Fold
  - Sorting
- Cool Overview
  - Syntax
  - Objects
  - Methods
  - Types



# This is my final day

- ... as your ... *companion* ... through Ocaml and Cool. After this we start the interpreter project.



# One-Slide Summary

- Functions and type inference are **polymorphic** and operate on more than one type (e.g., `List.length` works on int lists and string lists).
- **fold** is a powerful higher-order function (like a swiss-army knife or duct tape).
- **Cool** is a Java-like language with classes, methods, private fields, and inheritance.

# Pattern Matching (Error!)

- Simplifies Code (eliminates ifs, accessors)
  - `type btree = (* binary tree of strings *)`
  - `| Node of btree * string * btree`
  - `| Leaf of string`
  - `let rec height tree = match tree with`
  - `| Leaf _ -> 1`
  - `| Node(x,_,y) -> 1 + max (height x) (height y)`
  - `let rec mem tree elt = match tree with`
  - `| Leaf str | Node(_,str,_) -> str = elt`
  - `| Node(x,_,y) -> mem x elt || mem y elt`

bug?

# Pattern Matching (Error!)

- Simplifies Code (eliminates ifs, accessors)
  - `type btree =` (\* binary tree of strings \*)
  - | `Node` of `btree * string * btree`
  - | `Leaf` of `string`
  - `let rec bad tree elt = match tree with`
  - | `Leaf str | Node(_,str,_) -> str = elt`
  - | `Node(x,_,y) -> bad x elt || bad y elt`
  - `let rec mem tree elt = match tree with`
  - | `Leaf str | Node(_,str,_) when str = elt -> true`
  - | `Node(x,_,y) -> mem x elt || mem y elt`

# Pattern Matching Mistakes

- What if I forget a case?
  - `let rec is_odd x = match x with`
  - `| 0 -> false`
  - `| 2 -> false`
  - `| x when x > 2 -> is_odd (x-2)`
  - **Warning P: this pattern-matching is not exhaustive.**
  - **Here is an example of a value that is not matched: 1**

# Polymorphism

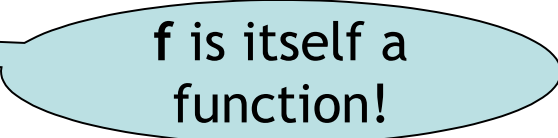
- Functions and type inference are polymorphic
  - Operate on more than one type
  - let rec length x = match x with
  - | [] -> 0
  - | hd :: tl -> 1 + length tl
  - val length :  $\alpha$  list -> int
  - length [1;2;3] = 3
  - length ["algol"; "smalltalk"; "ml"] = 3
  - length [1 ; "algol" ] = ?

$\alpha$  means "any one type"



# Higher-Order Functions

- Function are first-class values
  - Can be used whenever a value is expected
  - Notably, can be passed around
  - Closure captures the environment
  - `let rec map f lst = match lst with`
  - `| [] -> []`
  - `| hd :: tl -> f hd :: map f tl`
  - `val map : ( $\alpha$  ->  $\beta$ ) ->  $\alpha$  list ->  $\beta$  list`
  - `let offset = 10 in`
  - `let myfun x = x + offset in`
  - `val myfun : int -> int`
  - `map myfun [1;8;22] = [11;18;32]`
- Extremely powerful programming technique
  - General iterators
  - Implement abstraction



f is itself a function!

# The Story of Fold

- We've seen **length** and **map**

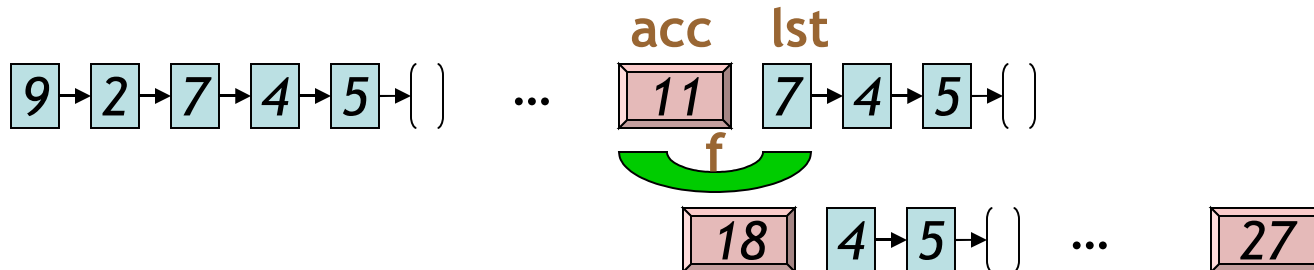
- We can also imagine ...

- **sum** [1; 5; 8 ] = 14
- **product** [1; 5; 8 ] = 40
- **and** [true; true; false ] = false
- **or** [true; true; false ] = true
- **filter** (fun x -> x>4) [1; 5; 8] = [5; 8]
- **reverse** [1; 5; 8] = [8; 5; 1]
- **mem** 5 [1; 5; 8] = true

- Can we build all of these?

# The House That Fold Built

- The **fold** operator comes from Recursion Theory (Kleene, 1952)
  - let rec fold f acc lst = match lst with
  - | [] -> acc
  - | hd :: tl -> fold f (f acc hd) tl
  - val fold :  $(\alpha \rightarrow \beta \rightarrow \alpha) \rightarrow \alpha \rightarrow \beta \text{ list} \rightarrow \alpha$
- Imagine we're summing a list (f = addition):



# It's Lego Time

- Let's build things out of Fold!
  - **length** lst = fold (fun acc elt -> acc + 1) 0 lst
  - **sum** lst = fold (fun acc elt -> acc + elt) 0 lst
  - **product** lst = fold (fun acc elt -> acc \* elt) 1 lst
  - **and** lst = fold (fun acc elt -> acc & elt) true lst
- How would we do **or**?
- How would we do **reverse**?



# Tougher Legos



- Examples:

- **reverse** lst = fold (fun acc e -> acc @ [e]) [] lst

- Note typing: (acc :  $\alpha$  list) (e :  $\alpha$ )

- **filter** keep\_it lst = fold (fun acc elt ->

- if keep\_it elt then elt :: acc else acc) [] lst

- **mem** wanted lst = fold (fun acc elt ->

- acc || wanted = elt) false lst

- Note typing: (acc : bool) (e :  $\alpha$ )

- How do we do **map**?

- Recall: map (fun x -> x +10) [1;2] = [11;12]

- Let's write it on the board ...

# Map From Fold

- let **map** myfun lst =
- fold (fun acc elt -> (myfun elt) :: acc) [] lst
  - Types: (myfun :  $\alpha \rightarrow \beta$ )
  - Types: (lst :  $\alpha$  list)
  - Types: (acc :  $\beta$  list)
  - Types: (elt :  $\alpha$ )
- How do we do **sort**?
  - (sort : ( $\alpha * \alpha \rightarrow \text{bool}$ ) ->  $\alpha$  list ->  $\alpha$  list)

*Do nothing which is of no use.*  
- Miyamoto Musashi, 1584-1645

# Sorting Examples

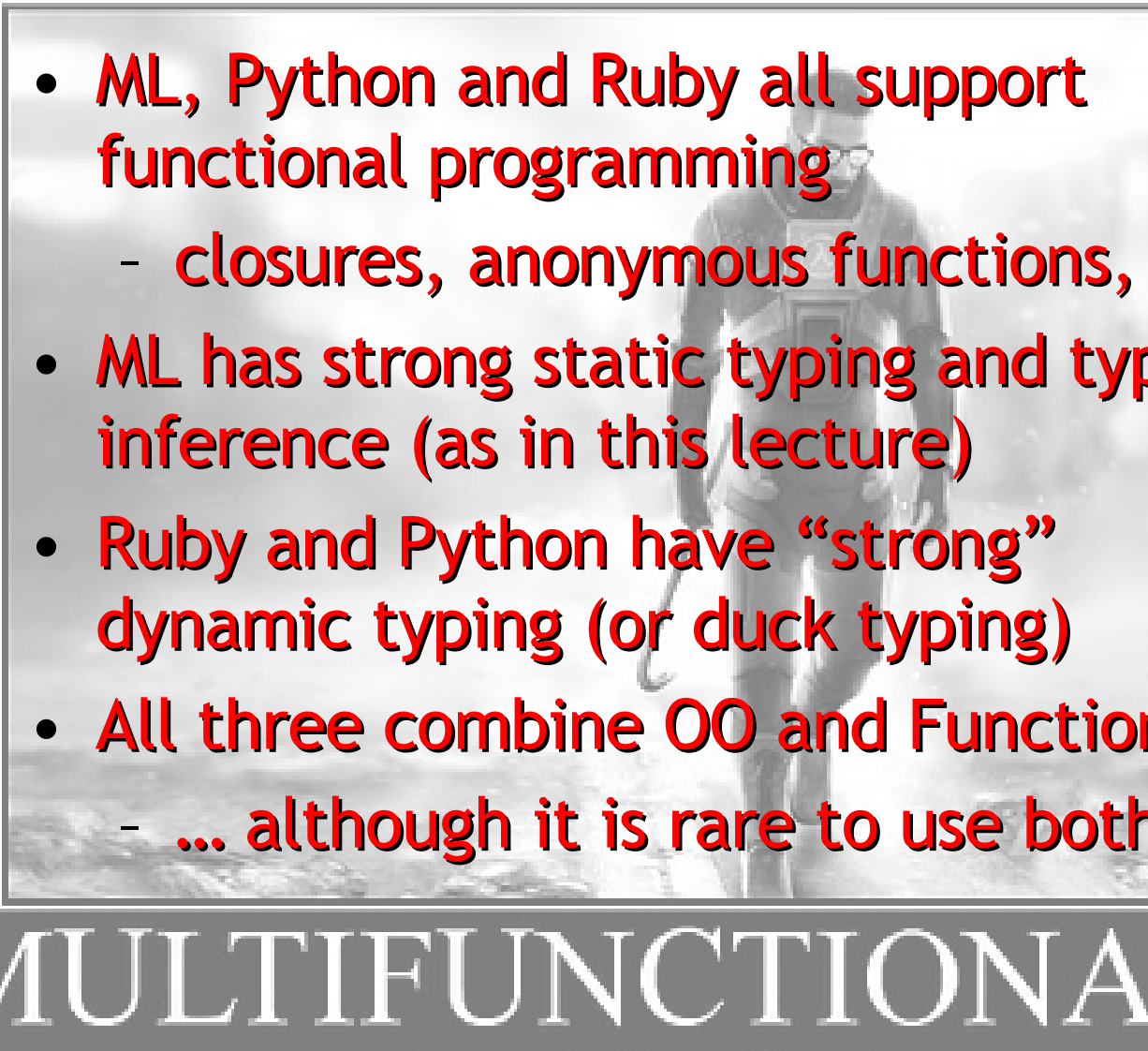
- `langs = [ "fortran"; "algol"; "c" ]`
- `courses = [ 216; 333; 415 ]`
- `sort (fun a b -> a < b) langs`
  - [ "algol"; "c"; "fortran" ]
- `sort (fun a b -> a > b) langs`
  - [ "fortran"; "c"; "algol" ]
- `sort (fun a b -> strlen a < strlen b) langs`
  - [ "c"; "algol"; "fortran" ]
- `sort (fun a b -> match is_odd a, is_odd b with`
  - | true, false -> true (\* odd numbers first \*)
  - | false, true -> false (\* even numbers last \*)
  - | \_, \_ -> a < b (\* otherwise ascending \*)) courses
  - [ 333 ; 415 ; 216 ]

*Java uses Inner Classes for this.*

# Partial Application and Currying

- let myadd x y = x + y
- **val myadd : int -> int -> int**
- myadd 3 5 = 8
- let addtwo = myadd 2
  - How do we know what this means? We use referential transparency! Basically, just substitute it in.
- **val addtwo : int -> int**
- addtwo 77 = 79
- Currying: “if you fix some arguments, you get a function of the remaining arguments”



- 
- A person wearing a backpack and glasses is walking through a field, looking at a laptop. The background is a soft-focus landscape with trees and a bright sky. The person is in the center of the frame, slightly to the right.
- **ML, Python and Ruby all support functional programming**
    - **closures, anonymous functions, etc.**
  - **ML has strong static typing and type inference (as in this lecture)**
  - **Ruby and Python have “strong” dynamic typing (or duck typing)**
  - **All three combine OO and Functional**
    - **... although it is rare to use both.**

# MULTIFUNCTIONALTY

One tool. One million uses.

# Cool Overview

- Classroom Object-Oriented Language
- Design to
  - Be implementable in one semester
  - Give a taste of implementing modern features
    - Abstraction
    - Static Typing
    - Inheritance
    - Memory management
    - And more ...
  - But many “grungy” things are left out

# A Simple Example

```
class Point {  
    x : Int <- 0;  
    y : Int <- 0;  
};
```

- Cool programs are sets of class definitions
  - A special **Main** class with a special method **main**
  - Like Java
- class = a collection of fields and methods
- Instances of a class are objects

# Cool Objects

```
class Point {  
    x : Int <- 0;  
    y : Int; (* use default value *)  
};
```

- The expression “**new Point**” creates a new object of class **Point**
- An object can be thought of as a record with a slot for each attribute (= field)

x	y
0	0

# Methods

```
class Point {  
  x : Int <- 0;  
  y : Int <- 0;  
  movePoint(newx : Int, newy : Int) : Point {  
    { x <- newx;  
      y <- newy;  
      self;  
    } -- close block expression  
  }; -- close method  
}; -- close class
```

- A class can also define methods for manipulating its attributes
- Methods refer to the current object using **self**

# Aside: Semicolons

```
class Point {  
  x : Int <- 0;  
  y : Int <- 0;  
  movePoint(newx : Int, newy : Int)  
    { x <- newx;  
      y <- newy;  
      self;  
    } -- close block  
}; -- close method  
}; -- close class
```

Yes, it's  
somewhat arbitrary.  
Still, don't get it wrong.



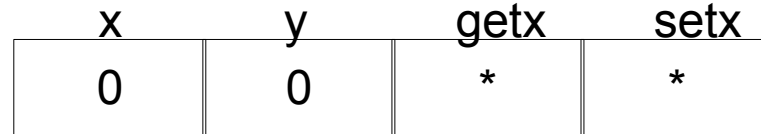
# Information Hiding

- Methods are **global**
- Attributes are **local** to a class
  - They can *only* be accessed by *that class's methods*

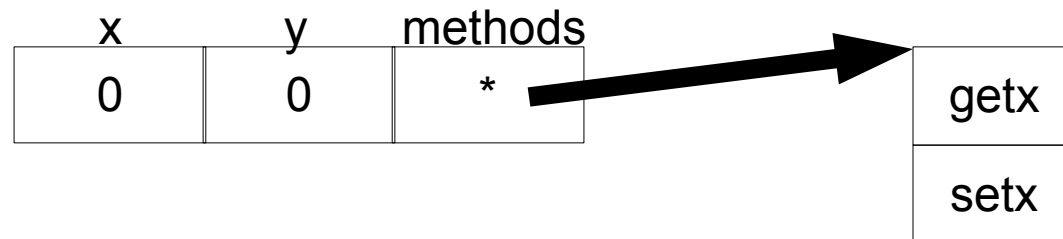
```
class Point {  
  x : Int <- 0;  
  y : Int <- 0;  
  getx () : Int { x } ;  
  setx (newx : Int) : Int { x <- newx };  
};
```

# Methods and Object Layout

- Each object knows how to access the code of its methods
- As if the object contains a slot pointing to the code



- In reality, implementations save space by sharing these pointers among instances of the same class





# Inheritance

- We can extend points to color points using **subclassing** => **class hierarchy**

```
class ColorPoint extends Point {  
  color : Int <- 0;  
  movePoint(newx:Int, newy:Int) : Point {  
    { color <- 0;  
      x <- newx; y <- newy;  
      self;  
    }  
  };  
};
```

Note references to fields x y –  
They're defined in Point!

x	y	color	movePoint
0	0	0	*

# Kool Types

- Every class is a **type**
- Base (built-in, predefined) classes:
  - **Int** for integers
  - **Bool** for booleans: **true**, **false**
  - **String** for strings
  - **Object** root of class hierarchy
- All variables must be declared
  - compiler infers types for expressions (like Java)



# Cool Type Checking

- **x : Point;**
- **x <- new ColorPoint;**
- ... is well-typed if **Point** is an ancestor of **ColorPoint** in the class hierarchy
  - Anywhere a **Point** is expected, a **ColorPoint** can be used (Liskov, ...)
- Rephrase: ... is well-typed if **ColorPoint** is a subtype of **Point**
- Type safety: a well-typed program *cannot* result in run-time type errors

# Method Invocation and Inheritance

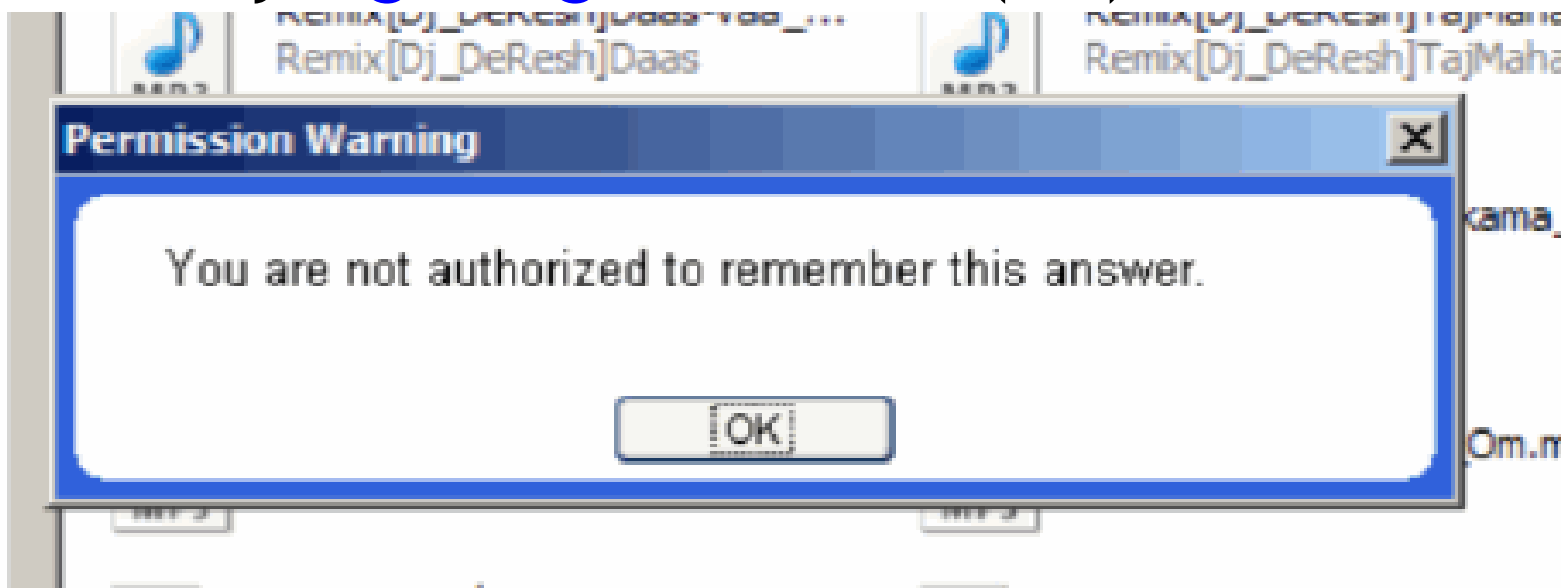
- Methods are invoked by (dynamic) dispatch
- Understanding dispatch in the presence of inheritance is a subtle aspect of OO
  - `p : Point;`
  - `p <- new ColorPoint;`
  - `p.movePoint(1,2);`
- `p` has static type `Point`
- `p` has dynamic type `ColorPoint`
- `p.movePoint` must invoke `ColorPoint` version

# Other Expressions

- Cool is an expression language (like Ocaml)
  - Every expression has a type and a value
  - Conditionals           if **E** then **E** else **E** fi
  - Loops                   while **E** loop **E** pool
  - Case/Switch           case **E** of **x** : **Type** => **E** ; ... esac
  - Assignment           **x** <- **E**
  - Primitive I/O         out\_string(**E**), in\_string(), ...
  - Arithmetic, Logic Operations, ...
- Missing: arrays, floats, interfaces, exceptions
  - Plus: you tell me!

# Cool Memory Management

- Memory is allocated every time “**new E**” executes
- Memory is deallocated automatically when an object is not reachable anymore
  - Done by a **garbage collector** (GC)



# Course Project

- A complete **interpreter**
  - Cool Source ==> Executed Program
  - No optimizations
  - Also no GC
- Split in 4 programming assignments (PAs)
- There is adequate time to complete assignments
  - But start early and follow directions
- PA2-4 ==> individual or teams (of max 2)

# Homework

- Wednesday: PA 0 due
- Thursday: Chapters 2.1 - 2.2
- Thursday: Dijkstra Paper
- 
- Bonus for getting this far: questions about fold are very popular on tests! If I say “write me a function that does foozle to a list”, you should be able to code it up with fold.