

# Profilers and Debuggers



# Introductory Material

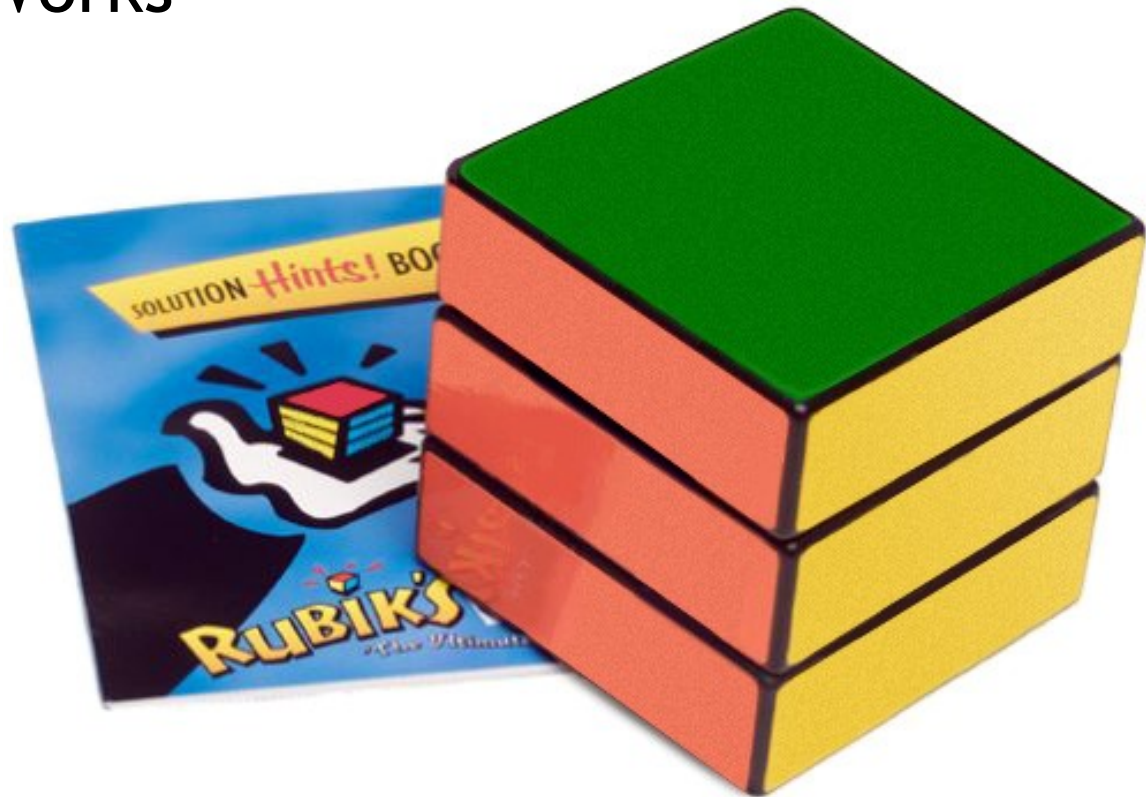
- First, who **doesn't** feel comfortable with assembly language?
  - You'll get to answer all the assembly questions. Yes, really.
- Lecture Style:
  - “Sit on the table” and pose questions. So, wake up!
- Lecture Goal:
  - After the lecture you'll think, “Wow, that was all really obvious. I could have done that.”

# One-Slide Summary

- A **debugger** helps to detect the source of a program error by **single-stepping** through the program and **inspecting** variable values.
- **Breakpoints** are the fundamental building block of debuggers. Breakpoints can be implemented with **signals** and **special OS support**.
- A **profiler** is a **performance** analysis tool that measures the frequency and **duration** of **function calls** as a program runs.
- Profilers can be **event-** or **sampling-based**.

# Lecture Outline

- Debugging
  - Signals
  - How Debuggers Works
  - Breakpoints
  - Advanced Tools
- Profiling
  - Event-based
  - Statistical



# What is a Debugger?

“A software tool that is used to detect the source of program or script errors, by performing step-by-step execution of application code and viewing the content of code variables.”

-Microsoft Developer Network

# Machine-Language Debugger

- Only concerned with **assembly code**
- Show instructions via **disassembly**
- Inspect the values of registers, memory
- Key Features (we'll explain all of them)
  - Attach to process
  - Single-stepping
  - Breakpoints
  - Conditional Breakpoints
  - Watchpoints

# Signals

- A **signal** is an **asynchronous** notification sent to a process about an event:
  - User pressed Ctrl-C (or did `kill %pid`)
    - Or asked the Windows Task Manager to terminate it
  - Exceptions (divide by zero, null pointer)
  - From the OS (**SIGPIPE**)
- You can install a **signal handler** - a procedure that will be executed when the signal occurs.
  - Signal handlers are vulnerable to **race conditions**.  
Why?

```
#include <stdio.h>
#include <signal.h>

int global = 11;

int my_handler() {
    printf("In signal handler, global = %d\n",
        global);
    exit(1);
}

void main() {
    int * pointer = NULL;

    signal(SIGSEGV, my_handler) ;

    global = 33;

    * pointer = 0;

    global = 55;

    printf("Outside, global = %d\n", global);
}
```

# Signal Example

- What does this program print?





# Attaching A Debugger

- Requires **operating system support**
- There is a special **system call** that allows one process to act as a debugger for a target
  - What are the **security** concerns?
- Once this is done, the debugger can basically “catch signals” delivered to the target
  - This isn't exactly what happens, but it's a good explanation ...

# Building a Debugger

```
#include <stdio.h>
#include <signal.h>

#define BREAKPOINT *(0)=0

int global = 11;

int debugger_signal_handler() {
    printf("debugger prompt: \n");
    // debugger code goes here!
}

void main() {
    signal(SIGSEGV, debugger_signal_handler);

    global = 33;

    BREAKPOINT;

    global = 55;

    printf("Outside, global = %d\n", global);
}
```

- We can then get breakpoints and interactive debugging
  - Attach to target
  - Set up signal handler
  - Add in exception-causing instructions
  - Inspect globals, etc.

# Reality

- We're not really changing the source code
- Instead, we modify the *assembly*
- We can't **insert** instructions
  - Because labels are already set at known constant offsets
- Instead we **change** them

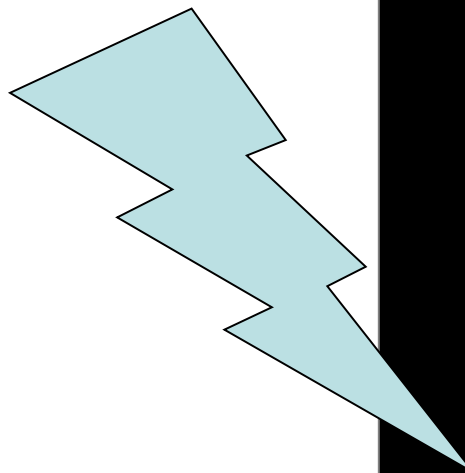


One of the class goals is to expose you to new languages: thus x86 ASM instead of COOL-ASM.

```
.file "example.c"
.globl _global
.data
.align 4
_global:
.long 11
.def __main
.section .rdata,"dr"
LC0:
.ascii "Outside, global = %d\12\0"
.text
.globl _main
.def _main
_main:
pushl %ebp
movl %esp, %ebp
subl $24, %esp
andl $-16, %esp
movl $0, %eax
addl $15, %eax
addl $15, %eax
shrl $4, %eax
sall $4, %eax
movl %eax, -4(%ebp)
movl -4(%ebp), %eax
call __alloca
call __main
movl $33, _global
movl $55, _global
movl _global, %eax
movl %eax, 4(%esp)
movl $LC0, (%esp)
call _printf
leave
ret
.def _printf
```

# Adding A Breakpoint

Add a breakpoint just after “global = 33;”



**Storage Cell:**

```
movl $55, _global
_main + 14
```

```
.file "example.c"
.globl _global
.data
.align 4
_global:
.long 11
.def __main
.section .rdata,"dr"
LC0:
.ascii "Outside, global = %d\12\0"
.text
.globl _main
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_main:
pushl %ebp
movl %esp, %ebp
subl $24, %esp
andl $-16, %esp
movl $0, %eax
addl $15, %eax
addl $15, %eax
shrl $4, %eax
sall $4, %eax
movl %eax, -4(%ebp)
movl -4(%ebp), %eax
call __alloca
call __main
movl $33, _global
movl $55, _global
movl _global, %eax
movl %eax, 4(%esp)
movl $LC0, (%esp)
call _printf
leave
ret
.def _printf
```

```
.file "example.c"
.globl _global
.data
.align 4
_global:
.long 11
.def __main
.section .rdata,"dr"
LC0:
.ascii "Outside, global = %d\12\0"
.text
.globl _main
.def _main
_main:
pushl %ebp
movl %esp, %ebp
subl $24, %esp
andl $-16, %esp
movl $0, %eax
addl $15, %eax
addl $15, %eax
shrl $4, %eax
sall $4, %eax
movl %eax, -4(%ebp)
movl -4(%ebp), %eax
call __alloca
call __main
movl $33, _global
movl $0, 0
movl _global, %eax
movl %eax, 4(%esp)
movl $LC0, (%esp)
call _printf
leave
ret
.def _printf
```

# Software Breakpoint Recipe

- Debugger has already attached and set up its signal handler
- User wants a breakpoint at instruction  $X$
- Store  $(X, \text{old\_instruction\_at\_}X)$
- Replace instruction at  $X$  with “\*0=0”
  - Pick something illegal that’s 1 byte long
- Signal handler replaces instruction at  $X$  with stored  $\text{old\_instruction\_at\_}X$
- Give user interactive debugging prompt

# Advanced Breakpoints

- Get register and local values by **walking the stack**
- Optimization: **hardware breakpoints**
  - Special register: if PC value = HBP register value, signal an exception
  - Faster than software, works on ROMs, only limited number of breakpoints, etc.
- Feature: **condition breakpoint**: “break at instruction *X* if **some\_variable = some\_value**”
- As before, but signal handler checks to see if **some\_variable = some\_value**
  - If so, present interactive debugging prompt
  - If not, return to program immediately
  - Is this fast or slow?

# Single-Stepping

- Debuggers allow you to advance through code on instruction at a time
- To implement this, put a breakpoint at the first instruction (= at program start)
- The “**single step**” or “**next**” interactive command is equal to:
  - Put a breakpoint at the next instruction
    - +1 for COOL-ASM, +4 bytes for RISC, +X bytes for CISC, etc.
  - Resume execution

# Watchpoints

- You want to know when a variable changes
- A **watchpoint** is like a breakpoint, but it stops execution whenever the value at location **L** changes, at any PC value
- How could we implement this?





# Watchpoint Implementation

- Software Watchpoints

- Put a breakpoint at *every instruction* (ouch!)
- Check the current value of **L** against a stored value
- If different, give interactive debugging prompt
- If not, set next breakpoint and continue (i.e., single-step)

- Hardware Watchpoints

- Special register holds **L**: if the value at address **L** ever changes, the CPU raises an exception

## Q: Movies (284 / 842)

- Name the movie described below and **either** the general scientific theory that Malcolm invokes or the ambushing cold-blooded killers. In this Oscar-winning 1993 Spielberg/Crichton extravaganza involving cloning and theme parks, Dr. Ian Malcolm correctly predicts that things will not turn out well.

Q: Advertising (799 / 842)

- Name the brand most associated with instant-print self-developing photographic film and cameras. The technology was invented in 1947 by corporation founder Edwin H. Land.

Q: Cartoons (671 / 842)

- Name all five main characters and the primary automobile from **Scooby Doo, Where Are You!**

# Real-World Languages

- This Northern European language boasts 5 million speakers (including Linus Torvalds). Its original writing system was devised in the 16<sup>th</sup> century from Swedish, German and Latin. Its eight vowels have powerful lexical and grammatical roles; doubled vowels do not become diphthongs.

# Source-Level Debugging

- What if we want to ...
  - Put a breakpoint at a *source-level* location (e.g., breakpoint at `main.c line 20`)
  - Single-step through *source-level* instructions (e.g., from `main.c:20` to `main.c:21`)
  - Inspect *source-level* variables (e.g., inspect `local_var`, not register AX)
- We'll need the compiler's help
- How can we do it?

# Debugging Information

- The compiler will emit tables
  - For every line in the program (e.g., main.c:20), what assembly instruction range does it map to?
  - For every line in the program, what variables are in scope *and where do they live* (registers, memory)?
- Put a breakpoint = table lookup
  - Put breakpoint at beginning of instruction range
- Single-step = table lookup
  - Put next breakpoint at end of instruction range +1
- Inspect value = table lookup
- Where do we put these tables?

These tables are conceptually similar to the class map or annotated AST.

# How Big Are Those Tables?

```
/* example.c */  
#include <stdio.h>  
#include <signal.h>  
  
int my_global_var = 11;  
  
void main() {  
  
    int my_local_var = 22;  
  
    my_local_var += my_global_var;  
  
    printf("Outside, my_local_var = %d\n", my_local_var);  
}
```

“gcc example.c”	9418 bytes
“gcc -g example.c”	23790 bytes



# Debugging vs. Optimizing

- We said: the compiler will emit tables
  - For every line in the program (e.g., main.c:20), what assembly instruction range does it map to?
  - For every line in the program, what variables are in scope and where do they live (registers, memory)?
- What can **go wrong** if we *optimize* the program?

# Replay Debugging

- Running and single-stepping are handy
- But wouldn't it be nice to go **back in time**?
- That is, from the current breakpoint, undo instructions in reverse order

- Intuition: functional + single assignment

`x = 11;`

`x = x + 22;`

`breakpoint ;`

`x = x + 33;`

`print x`



`let x0 = 11 in`

`let x1 = x0 + 22 in`

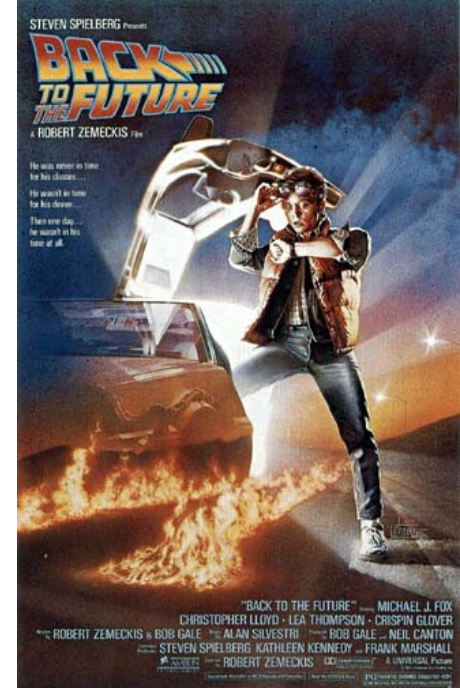
`breakpoint ;`

`let x2 = x1 + 33 in`

`print x`

# Time Travel

- **Store the state** at various times
  - time  $t=0$  at program start
  - time  $t=88$  after 88 instructions
  - ... *why does this work?*
- When the user asks you to go back one step, you actually *go back to the last stored state* and run the program forward again with a breakpoint
  - e.g., to go back from  $t=150$ , put breakpoint at instruction 149 and re-run from  $t=88$ 's state
- **ocamldebug** has this power - try it!



# Valgrind

- **Valgrind** is a suite of free tools for debugging and profiling
  - Finds **memory errors**, profiles cache times, call graphs, profiles heap space
- It does so via **dynamic binary translation**
  - Fancy words for “it is an interpreter”
  - No need to modify, recompile or relink
  - Works with any language
- Can attach gdb to your process, etc.
- Problem: slowdown of 5x-100x
  - Rational Purify (commercial) is similar
  - Check out Strata by Jack Davidson



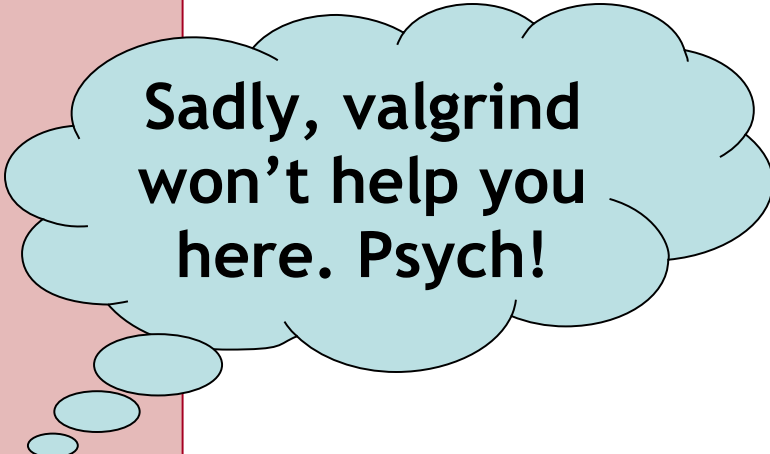
# Valgrind Example

```
int main() {  
    int some_var = 55;  
    int array[10];  
    int i;  
    for (i=0;i<=10;i++)  
        array[i] = i;  
    printf("some_var = %d\n",  
        some_var);  
}
```

*What's the output?*

# Valgrind Example

```
int main() {  
    int some_var = 55;  
    int array[10];  
    int i;  
    for (i=0;i<=10;i++)  
        array[i] = i;  
    printf("some_var = %d\n",  
        some_var);  
}
```



Sadly, valgrind  
won't help you  
here. Psych!

```
[weimer@weimer-laptop ~]$ ./a.out  
some_var = 10
```

# DDD

- Gnu Data Display Debugger
  - Similar in spirit to Visual Studio's built-in debugger
  - But for gdb, the Java debugger, the perl debugger, the python debugger, etc.
- How does this work? You tell me!

DDD: /public/source/programming/ddd-3.2/ddd/cxxtest.C

File Edit View Program Commands Status Source Data Help

0: list->self

1: list  
(List \*) 0x804df80

value = 85  
self = 0x804df80  
next = 0x804df90

value = 86  
self = 0x804df90  
next = 0x804df90

```
list->next = new List(a_global + start++);  
list->next->next = new List(a_global + start++);  
list->next->next->next = list;
```

(void) list; // Display this

(List \*) 0x804df80

delete list; // none money  
delete list->next;  
delete list;

// Test  
void lis  
{  
list  
}

//  
void ref  
{  
date  
dele  
date\_per

DDD Tip of the Day #5

If you made a mistake, try **Edit→Undo**. This will undo the most recent debugger command and redisplay the previous program state.

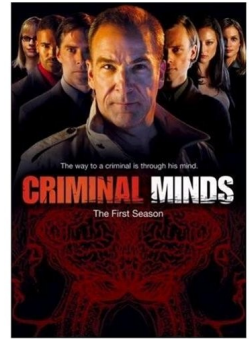
Close Prev Tip Next Tip

(gdb) graph display \*(list->next->next->self) dependent on 4  
(gdb) ↓

list = (List \*) 0x804df80



# Profiling



- A **profiler** is a performance analysis tool that measures the frequency and duration of function calls as a program runs.
- **Flat profile**
  - Computes the average call times for functions but does not break times down based on context
- **Call-Graph profile**
  - Computes call times for functions and also the call-chains involved

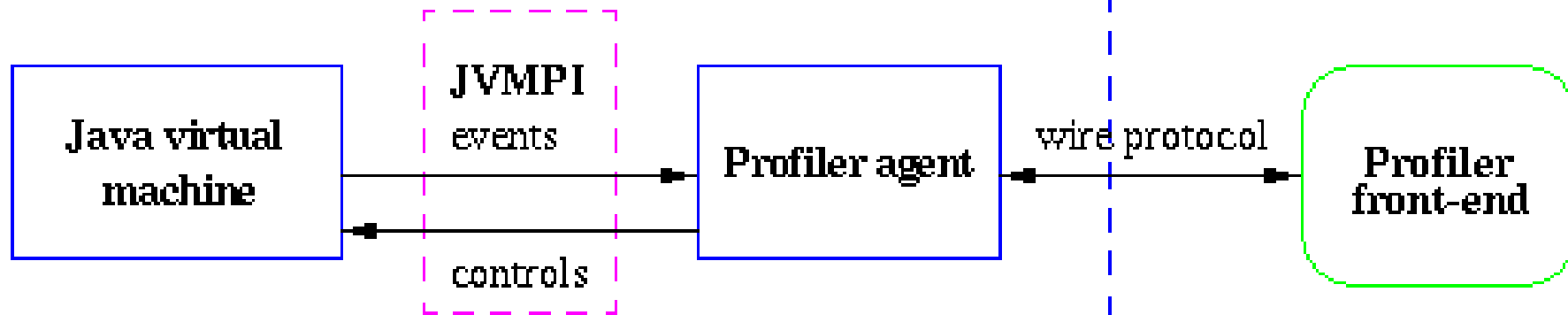


# Event-Based Profiling

- **Interpreted languages** provide special hooks for profiling
  - Java: JVM-Profile Interface, JVM API
  - Python: `sys.set_profile()` module
  - Ruby: `profile.rb`, etc.
- You **register a function** that will get called whenever the target program calls a method, loads a class, allocates an object, etc.
  - You could do this for PA5: count the number of object allocations, etc.
  - (And we do some profiling for you in PA7.)

# JVM Profiling Interface

- VM notifies profiler agent of various **events** (heap allocation, thread start, method invocation, etc.)
- Profiler agent issues control commands to the JVM and communicates with a GUI



Java virtual machine process

Profiler process

# Statistical Profiling



- You can arrange for the operating system to send you a **signal** (just like before) every X seconds (see `alarm(2)`)
- In the **signal handler** you determine the value of the target **program counter**
  - And append it to a growing list file
  - This is called **sampling**
- Later, you use that debug information table to map the PC values to procedure names
  - Sum up to get amount of time in each procedure

# Sampling Analysis

- Advantages
  - Simple and cheap - the **instrumentation** is unlikely to disturb the program too much
  - No big slowdown
- Disadvantages
  - Can completely miss periodic behavior (e.g., you sample every  $k$  seconds but do a network send at times  $0.5 + nk$  seconds)
  - **High error rate**: if a value is  $n$  times the sampling period, the expected error in it is  $\sqrt{n}$  sampling periods
- Read the `gprof` paper for midterm2

# While Derivation On The Board?

- If we have time, let's do this together ...
- $E = [ x \rightarrow a ]$
- $S = [ a \rightarrow 0 ]$
- $S' = [ a \rightarrow 1 ]$

**while x < 1 loop x <- x + 1 pool**

# Homework

- **Midterm 2** - Mon April 21<sup>st</sup> In Class
  - Covers Lectures “Code Generation” to “Linking, Loading and Shared Libraries” (i.e., everything after Midterm 1) plus each WA and PA done during that time
  - Everything *after* Earley parsing
- **Midterm 2 Review Session?**
  - Post of the forum!