

Written Assignment 5

This assignment asks you to prepare written answers to questions on code layout and operational semantics. Each of the questions has a short answer. You may discuss this assignment with other students and work on the problems together. However, your write-up should be your own individual work.

Please print your name and email address on your homework!

We need this information so that we can give you credit for the assignment and so that we can return it to you.

1. Consider the following piece of code:

```
fact (n : Int) : Int {
  if n > 0 then n*fact (n-1) else 1 fi
};
```

Draw the tree of activation records for a call to `fact (4)`. Include at least the stack pointers, frame pointers, return values, and parameters.

2. Consider these six operational semantics rules:

$$\begin{array}{l}
 (1) \frac{so, E, S \vdash e_1 : Bool(false), S_1}{so, E, S \vdash \text{while } e_1 \text{ loop } e_2 \text{ pool} : void, S_1} \qquad (4) \frac{E(id) = l_{id} \quad S(l_{id}) = v}{so, E, S \vdash id : v, S} \\
 \frac{so, E, S \vdash e_1 : Bool(true), S_1 \quad so, E, S_1 \vdash e_2 : v, S_2}{so, e, S_2 \vdash \text{while } e_1 \text{ loop } e_2 \text{ pool} : void, S_3} \quad (5) \frac{so, E, S \vdash e : v, S_1 \quad E(id) = l_{id} \quad S_2 = S_1[v/l_{id}]}{so, E, S \vdash id \leftarrow e : v, S_2} \\
 (3) \frac{so, E, S \vdash e_1 : v_1, S_1 \quad l_{new} = newloc(S_1) \quad so, E[l_{new}/id], S_1[v_1/l_{new}] \vdash e_2 : v_2, S_2}{so, E, S \vdash \text{let } id : T \leftarrow e_1 \text{ in } e_2 : v_2, S_2} \\
 (6) \frac{so, E, S \vdash e_1 : Int(n_1), S_1 \quad so, E, S_1 \vdash e_2 : Int(n_2), S_2 \quad v = \begin{cases} Bool(true) & \text{if } n_1 < n_2 \\ Bool(false) & \text{if } n_1 \geq n_2 \end{cases}}{so, E, S \vdash e_1 < e_2 : v, S_2}
 \end{array}$$

Use these rules to construct a derivation for the following piece of code:

```

let x : Int <- 2 in
while 1 < x loop
  x <- x - 1
pool

```

You may assume reasonable axioms, e.g. it is always true that $so, E, S \vdash 2 - 1 : Int(1), S$. Start your derivation using the `let` rule (3) as follows:

$$\frac{\frac{so, E, S \vdash 2 : Int(2), S}{so, E[l_{new}/x], S[Int(2)/l_{new}] \vdash \text{while } 1 < x \text{ loop } x \leftarrow x - 1 \text{ pool} : void, S_{final}}{\dots}}{so, E, S \vdash \text{let } x : Int \leftarrow 2 \text{ in while } 1 < x \text{ loop } x \leftarrow x - 1 \text{ pool} : void, S_{final}} \quad (3)$$

Note that you only need to expand hypotheses that need to be proved (i.e. those containing \vdash).

3. Suppose we wanted to add arrays to COOL, using the following syntax:

```

let a:T[e1] in e2   Create an array a with size e1 of Ts, usable in e2
a[e1] <- e2       Assign e2 to element e1 in a
a[e]                 Get element e of a

```

Write the operational semantics for these three syntactic constructs. You may find it helpful to think of an array of type $T[n]$ as an object with n attributes of type T .

4. The operational semantics for COOL's `while` expression show that result of evaluating such an expression is always `void`.

However, we could have used the following alternative semantics:

- If the loop body executes at least once, the result of the `while` expression is the result from the *last* iteration of the loop body.
- If the loop body never executes (i.e., the condition is false the first time it is evaluated), then the result of the `while` expression is `void`.

For example, consider the following expression:

```

while (x < 10) loop x <- x+1 pool

```

The result of this expression would be 10 if, initially, $x < 10$ or `void` if $x \geq 10$. Write new operational rules for the `while` construct that formalize these alternative semantics.