

Understanding Automatically- Generated Patches Through Symbolic Invariant Differences

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The Problem

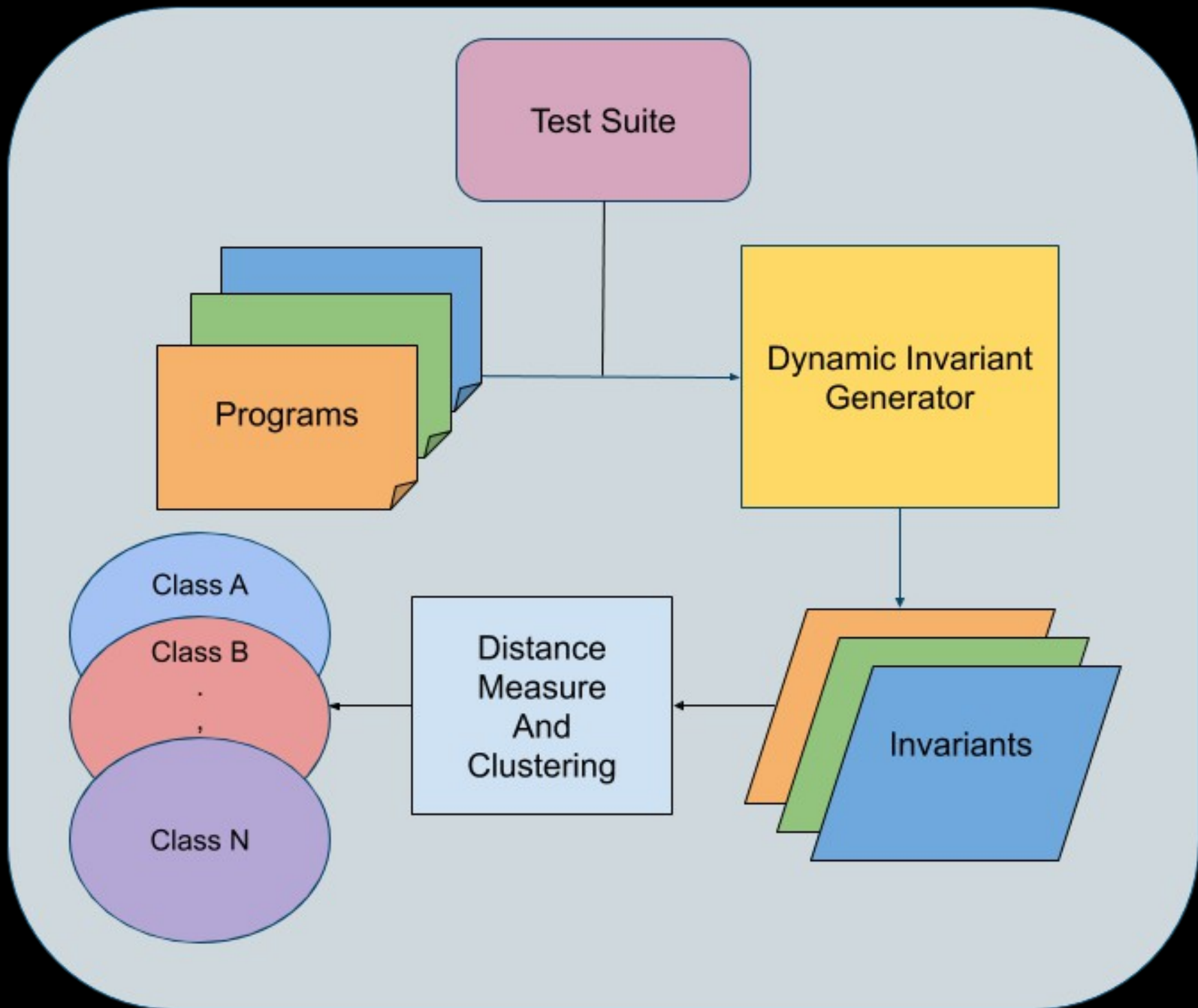
- **Automated program repair** may reduce software maintenance costs
 - Given a program and evidence of a bug, produce patches that fix that bug
 - SapFix, Angelix, Hercules, Prophet, Darjeeling, ...
- A **plausible patch** passes local tests but *may or may not* be acceptable to developers
 - Assessing plausible patches takes time and effort
 - Can we reduce that manual analysis time?

Patch Quality

- Many quality properties influence human decisions to adopt patches
 - Readability, maintainability, trust, style, ...
- In addition, there are functional correctness concerns related to **overfitting**
- Repair algorithms may incorporate techniques to produce more acceptable patches
 - (e.g., templates, restricted operators, consolidation, etc.)

Patch Assessment

- Ultimately, generate-and-validate program repair may produce **dozens of syntactically-unique patches** for the same defect
- We propose to **reduce this inspection burden**
 - Characterize patches by their sets of formal invariants (i.e., their behavior)
 - Calculate a distance metric on invariant sets
 - Cluster invariant sets (and thus patches) into equivalence classes
 - Only inspect one patch of each equivalence class



Comparing Invariant Sets

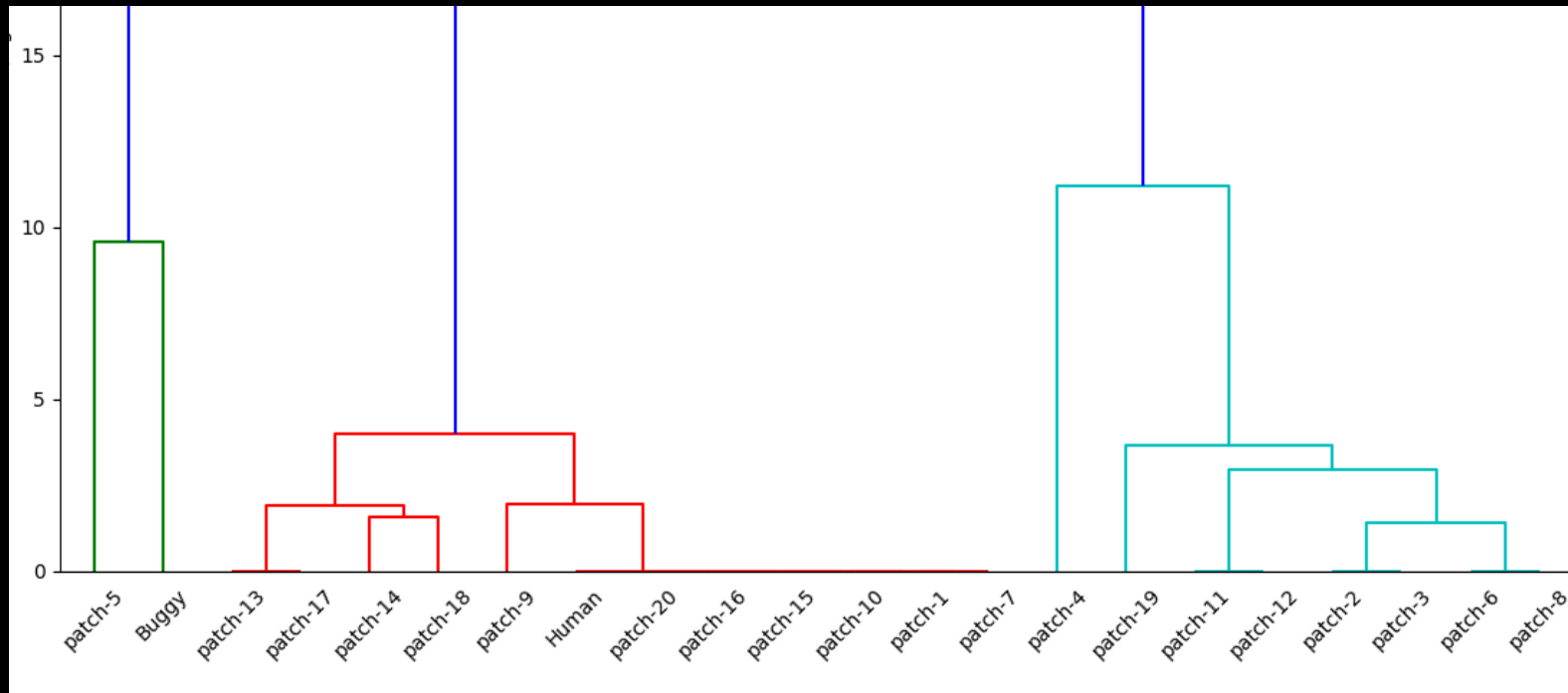
- Relaxes standard set difference from requiring equivalence to requiring logical implication
- Given programs A and B, tests T and invariant sets I_A and I_B
- We define the **implication distance** to be the cardinality of the subset of invariants in I_B that are *not implied* by any invariant in I_A
 - This definition admits hierarchical clustering
 - Optimization: consider only minterms from I_A

Efficient Invariant Comparison

- We also consider a more **syntactic** notion of distance on invariant sets
- We map syntactically-identical invariants to the same logical alphabet symbol
 - “ $X=2$ ” is A, “ $X=2$ ” is A, “ $X=1+1$ ” is B, etc.
- And then calculate the **Levenshtein edit distance on the induced strings**
 - Efficient polytime computation (cf. Z3)

Results & Conclusion

- Applied to 7 Defects4J and 5 ManyBugs bugs
 - 20-50 patches each from multiple tools



- **Reduces manual inspection burden by 40-50%**
 - Fast string-based distance has **95% accuracy**