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# Towards Autonomous Topological Place Detection Using the Extended Voronoi Graph

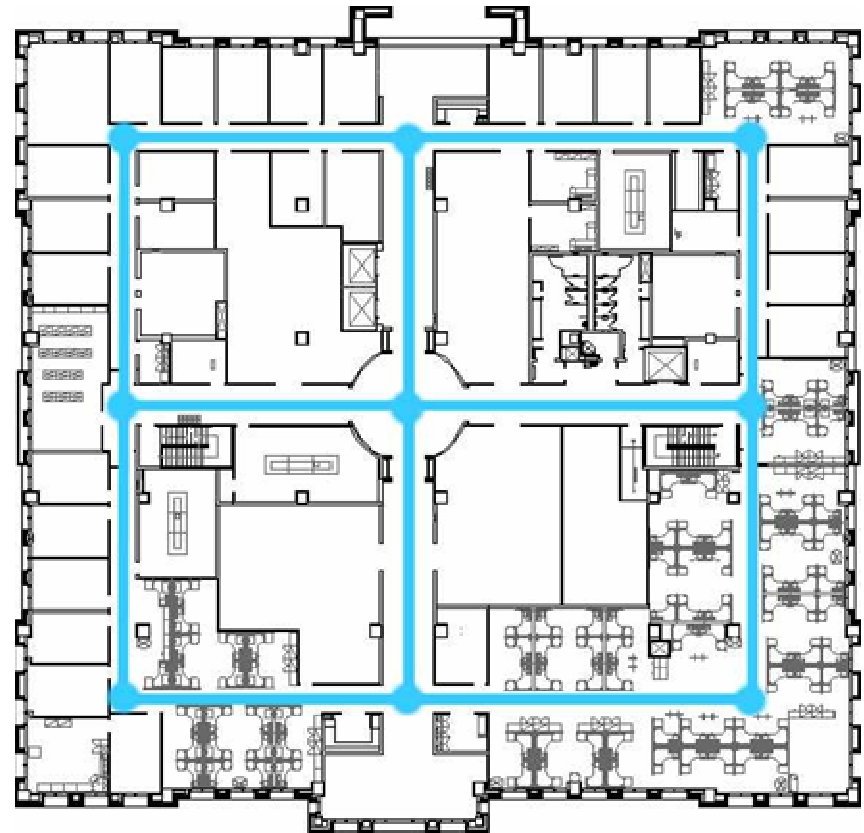
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# Problem Assumptions

- This work focuses on dead-ends, intersections, and doorways.
  - Smallest, atomic places sufficient for topological modeling.
  - Psychologically motivated



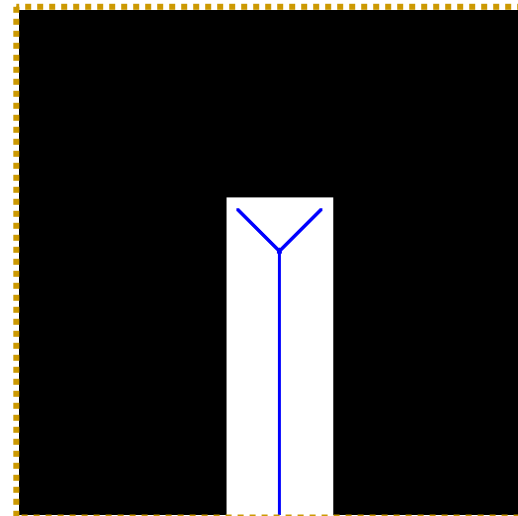
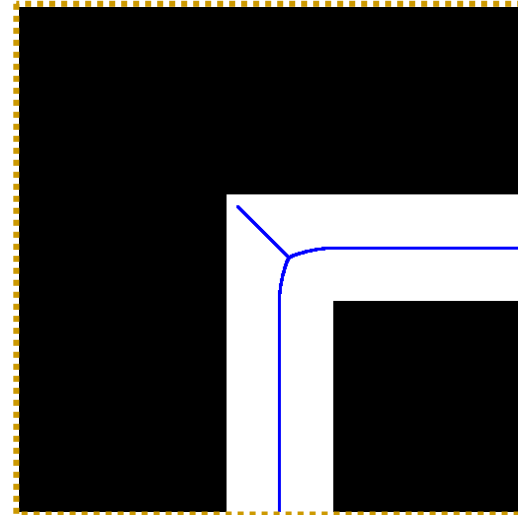
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# Previous Solutions

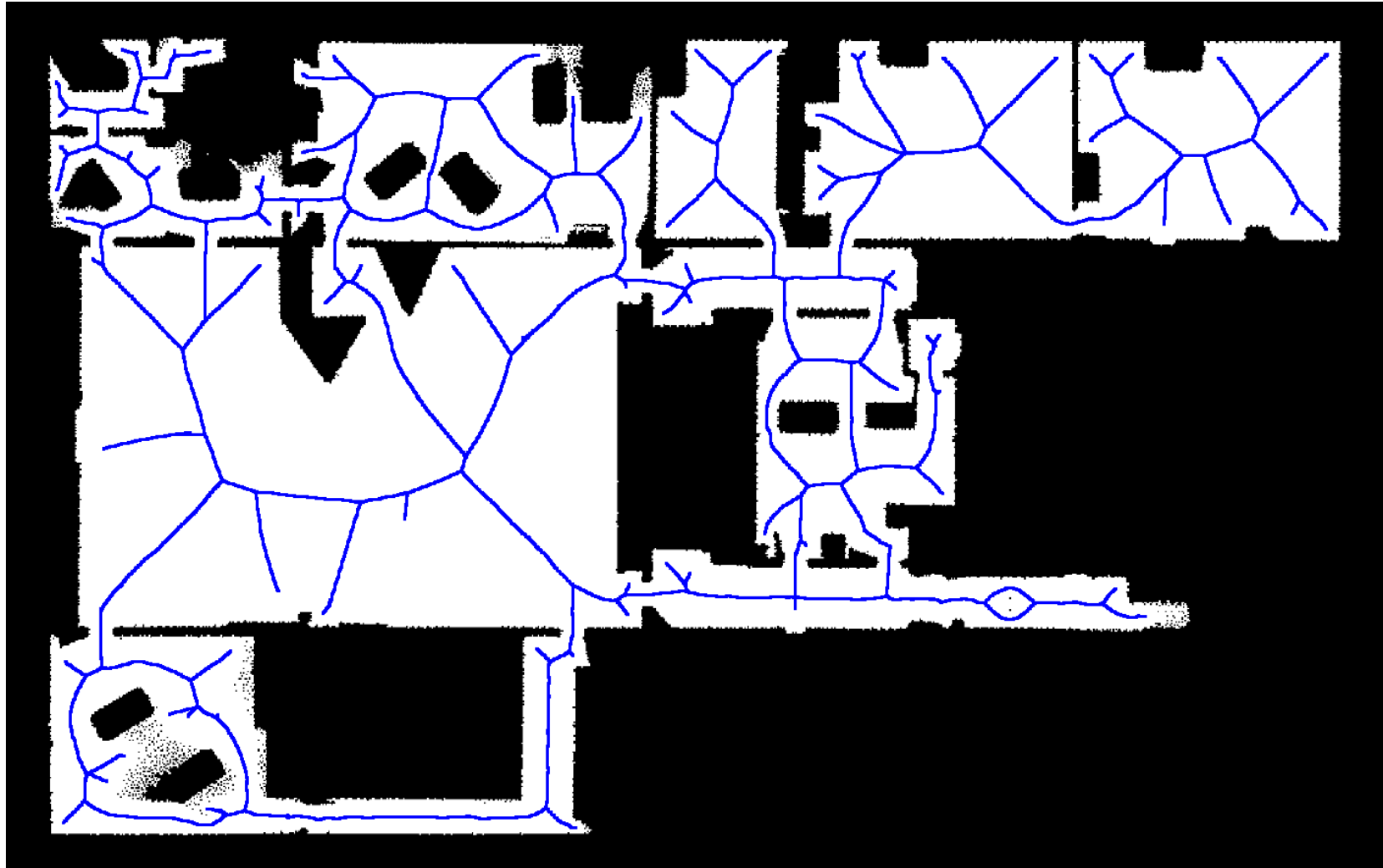
- Ad-hoc detection of simple intersections
  - Hand-coded
  - Environment dependent
  - Brittle
- Voronoi graph defines paths
  - Principled approach
    - Well-defined in enclosed environments
    - Work even in unconventional / non-perpendicular intersections
  - Experimentally validated in a number of environments
  - Choset and Nagatani, Trans. Robotics & Automation, 2001

# Places at Voronoi Graph Junctions

- Voronoi graph – 1D set of points equidistant to the  $N$  (or more) closest obstacles in  $N$  dimensions.
- Usually:
  - Wheeled mobile robots
  - Planar range-sensors
  - 2D occupancy grid
- Places – branching points in the Voronoi graph
  - junctions



# Voronoi Graph of a Global Metrical Map



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# Voronoi Problems

- Detecting Voronoi junctions has demonstrated some success for autonomous place detection.
- However, there are well known problems:
  - Unreliable
    - Spurious detections
    - Missed detections
    - Multiple places in complex intersections
  - Environment dependent
    - Only works in “enclosed” regions (at least N nearby obstacles)

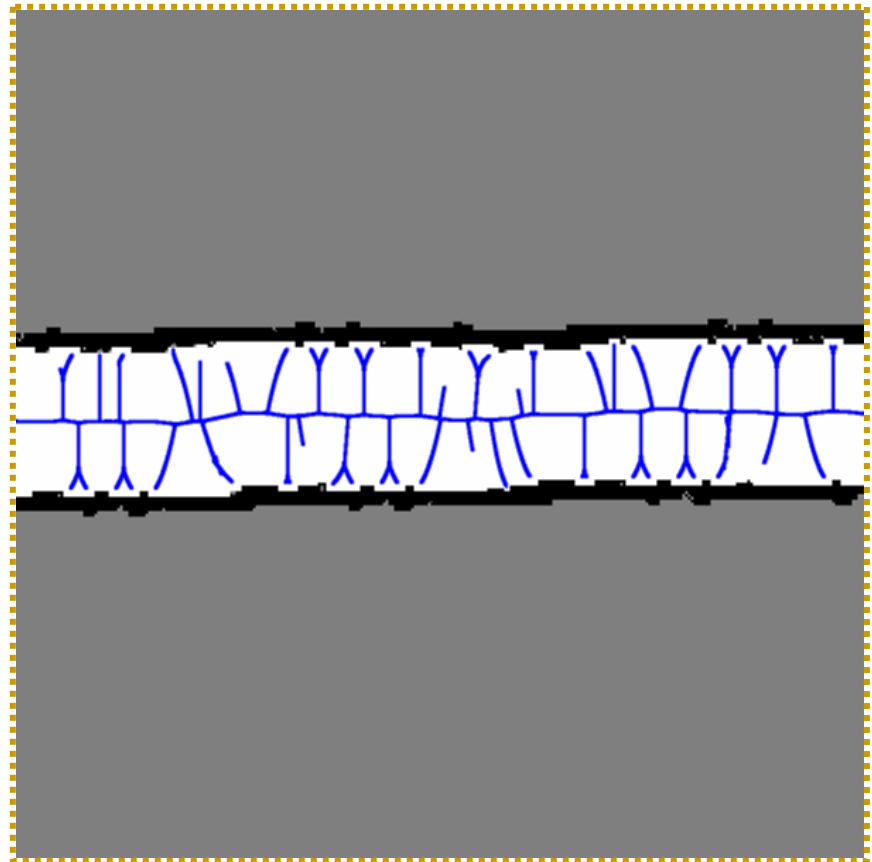
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# Our Contributions

1. Provide a principled approach to pruning the Voronoi graph.
  - Reduces spurious Voronoi branches (i.e. paths)
2. Provide a simple extension to the standard Voronoi graph definition.
  - Allows non-enclosed (e.g. outdoor) environments
3. Provide a definition of intersections that is better than Voronoi junctions.
  - Reduces missed places
  - Aggregates regions at complex intersections into places

# Pruning the Voronoi Graph

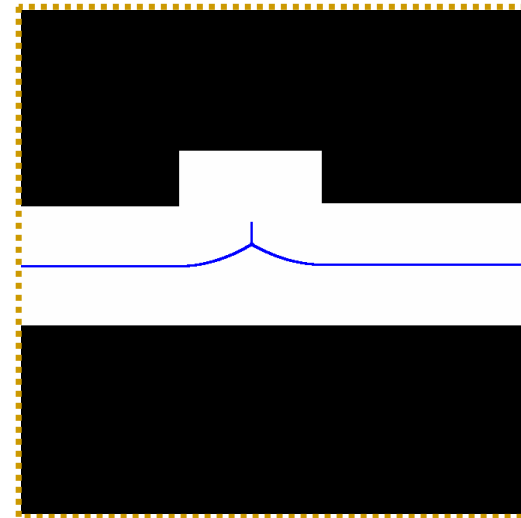
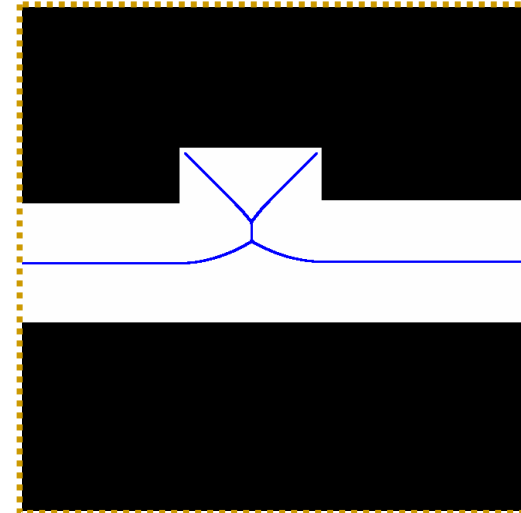
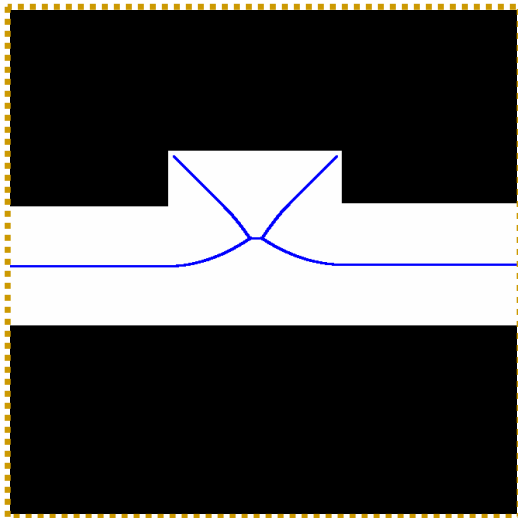
- Pruning is necessary to remove spurious branches.
- Choset's method
  - Remove terminal branches that only touch 2 obstacles.





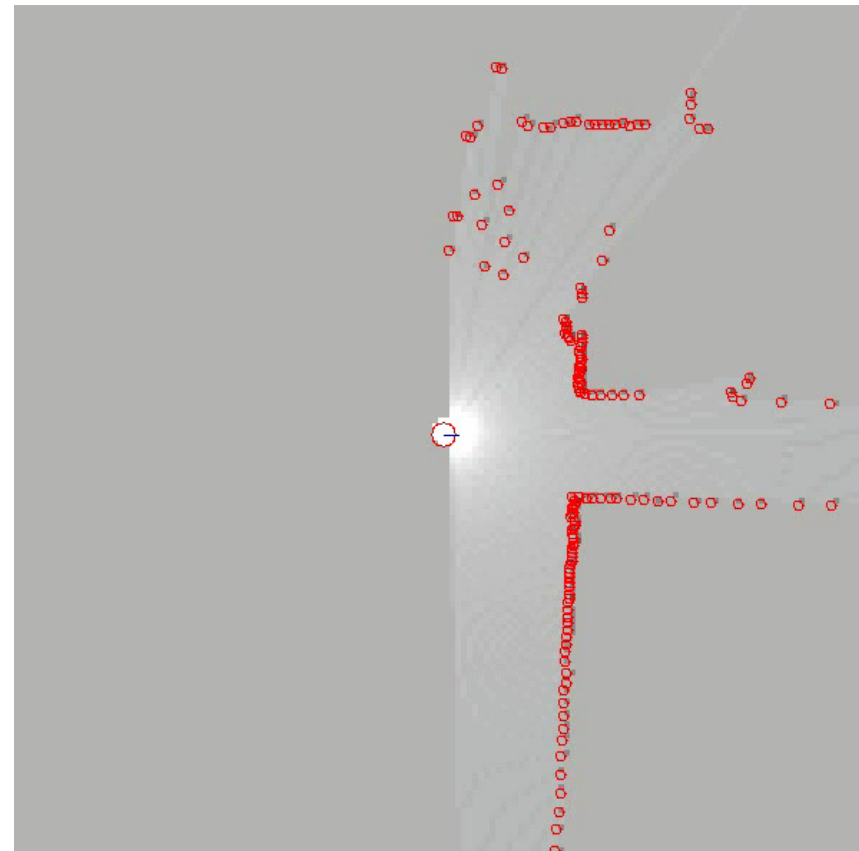
# Pruning Failures

- Fixed-depth pruning fails for hierarchical branching.



# Our Solution to Pruning (Background)

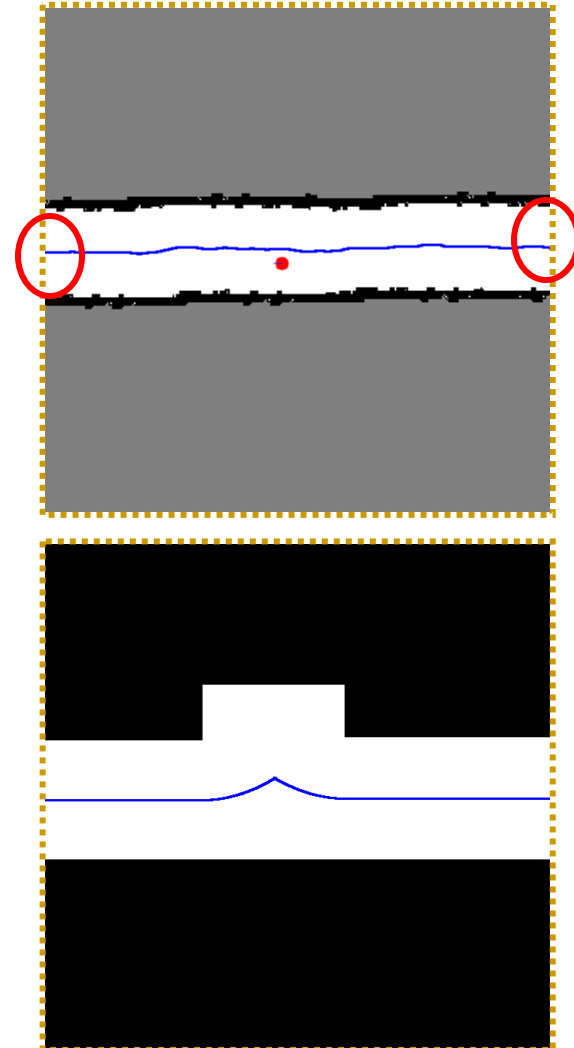
- There is always a fixed size, scrolling occupancy grid centered near the robot.
- (From previous work)
  - Small-scale space – local metrical model
  - Large-scale space – global topological model
  - If desired – global metrical model built using the topological skeleton



Video

# Our Solution to Pruning

- Define **exits** as points in the Voronoi graph that:
  - touch the grid edge
  - touch “unknown” space (gray cells) in the grid
- Find the minimum spanning tree that connects the exits.
  - Dead-ends are a simple special case.
- The size of the scrolling, local occupancy grid will determine the size of relevant places and paths.



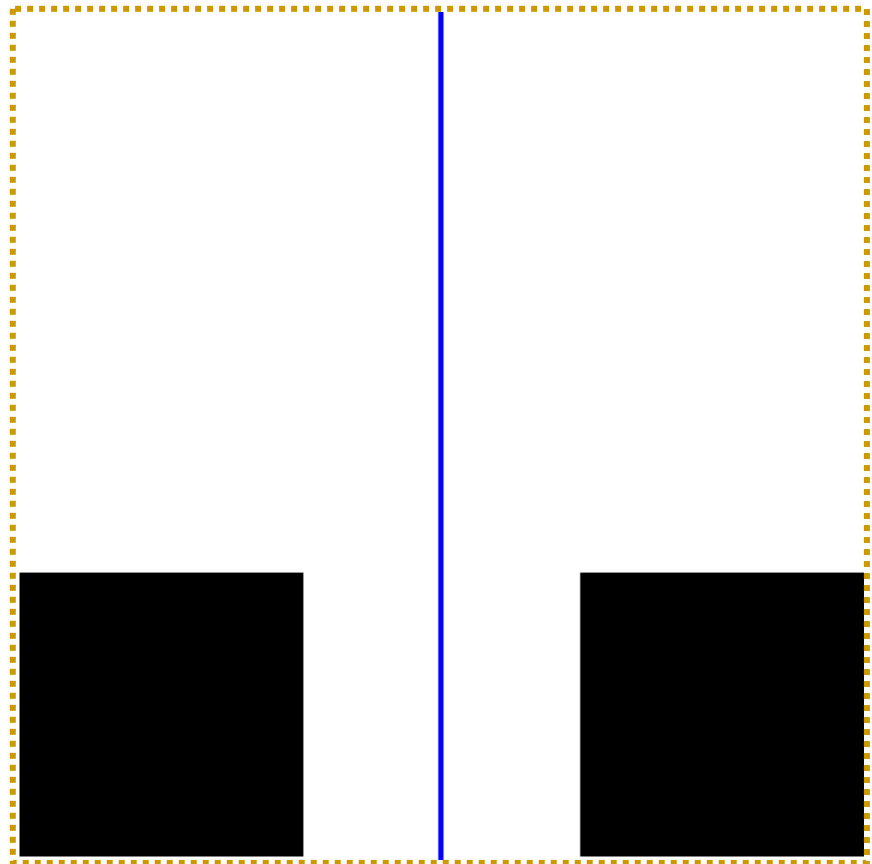
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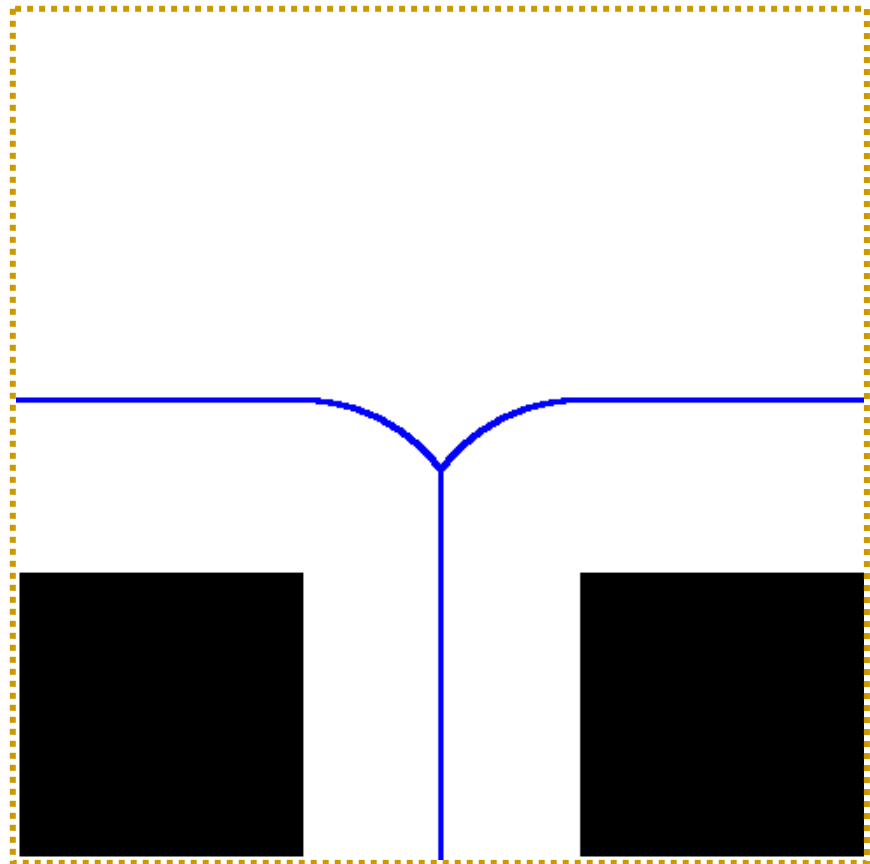
# Extending the Voronoi Graph

- Voronoi graphs require  $N$  or more nearby obstacles.
  - In our case  $N=2$ .
- 1. This limits the robot to exploring enclosed environments.
- 2. If following the Voronoi graph, the robot can become lost.

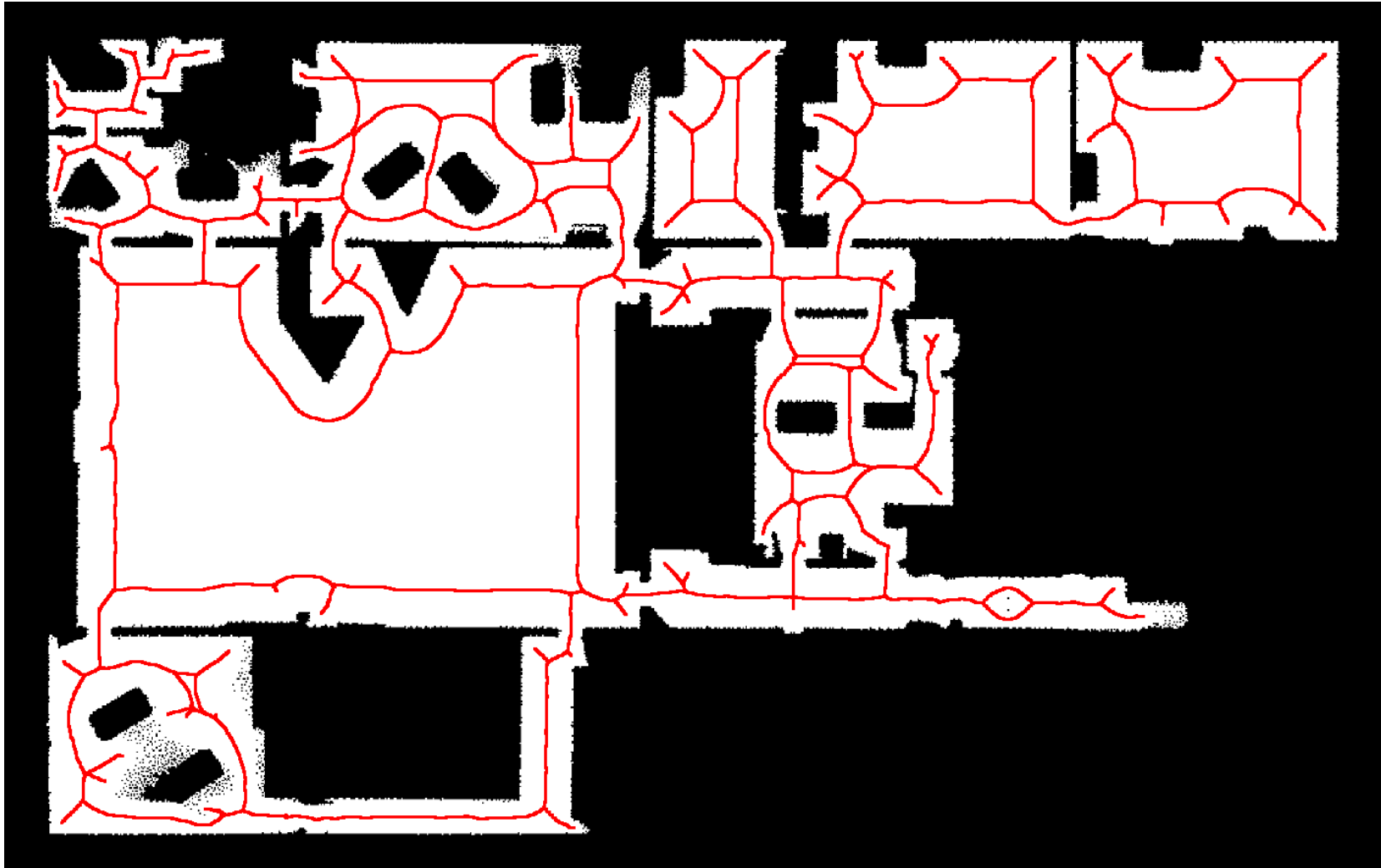


# Extended Voronoi Graph

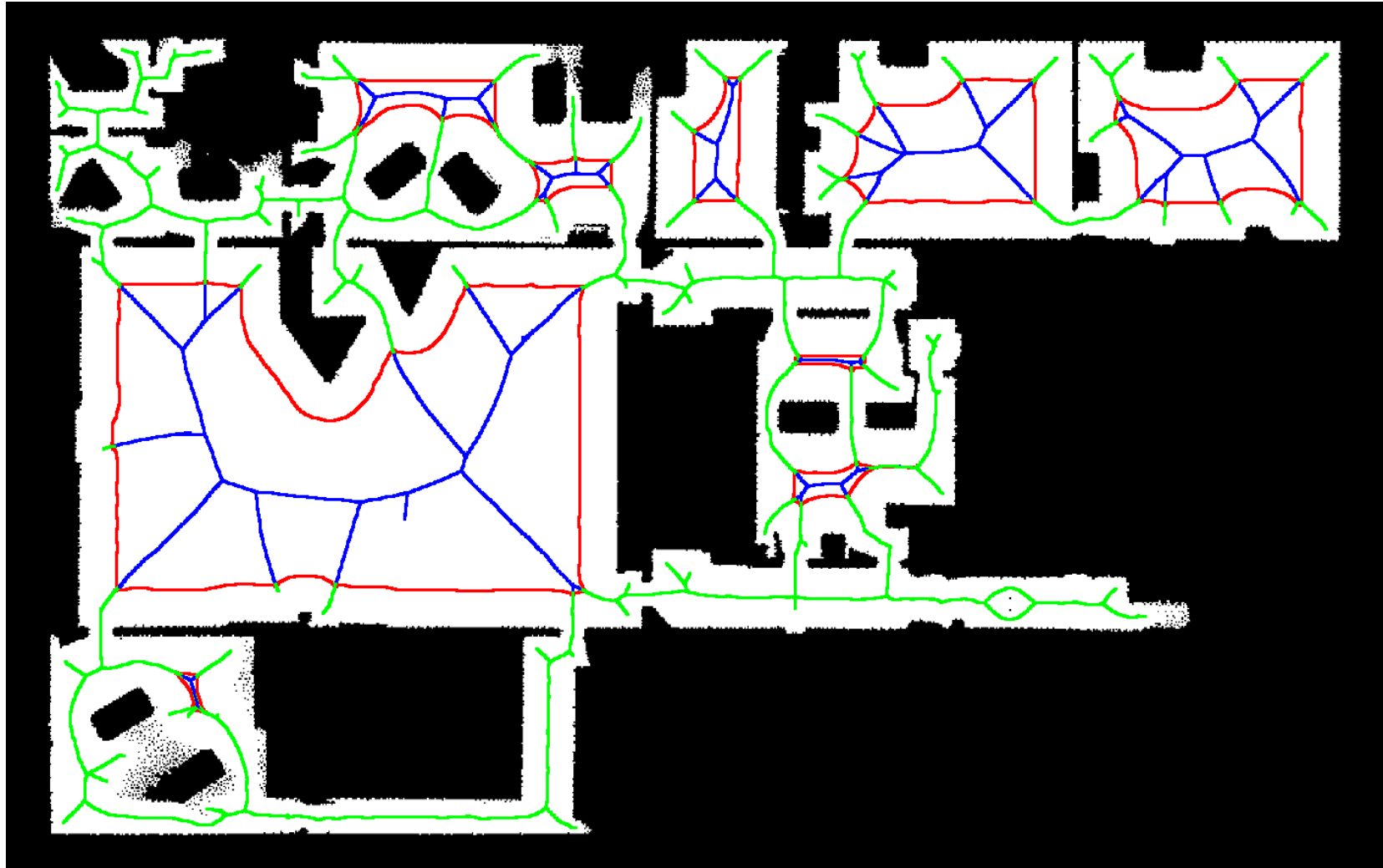
- UNION of:
  - Subset of the regular Voronoi graph, where  $r < d < S$ .
    - $d$  is distance to nearest obstacle.
    - $r$  is based on the robot's physical body.
    - $S$  is a maximum distance threshold
  - The set of points where  $d = S$ .
    - Provides “coastal navigation” when necessary.



# Extended Voronoi Graph



# Voronoi Graph Comparison





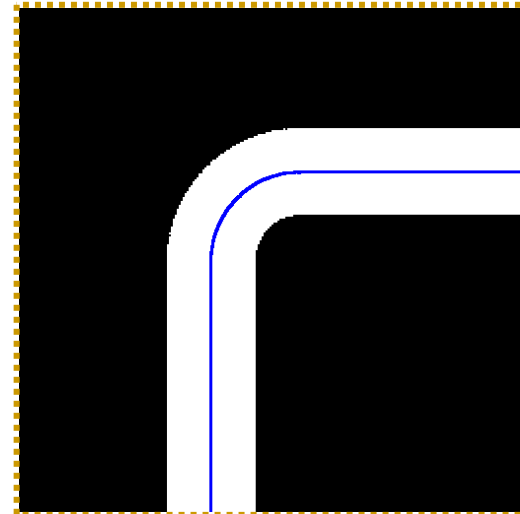
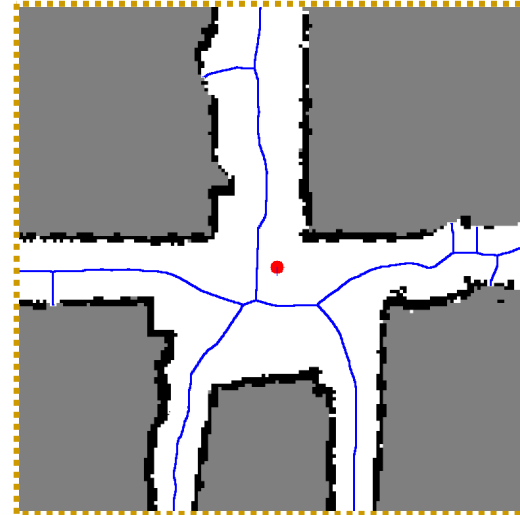
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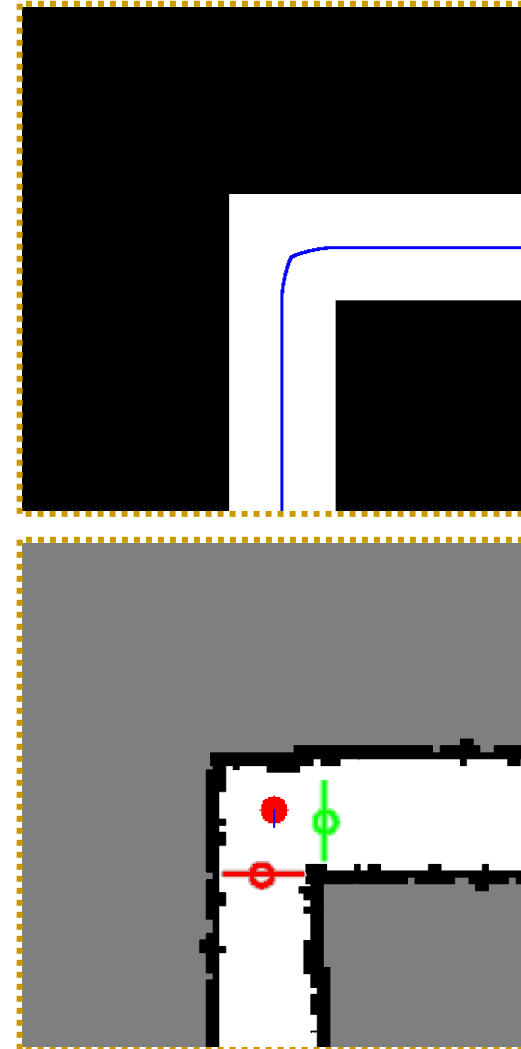
# Problems with Places at Junctions

- Junctions are highly susceptible to environmental / sensor noise.
- Some places have multiple junctions.
- Some places do not have junctions.



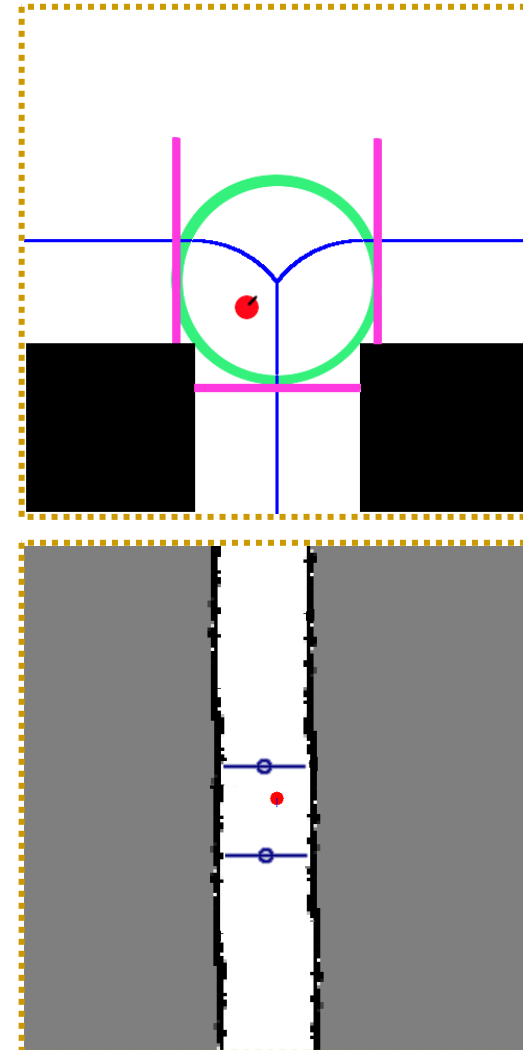
# Problems with Paths as Branches

- Using branches to define paths results in multiple places in complex environments.
- Sometimes there are no places in simple environments.
- Our approach uses *gateways* to find paths in the local metrical model.

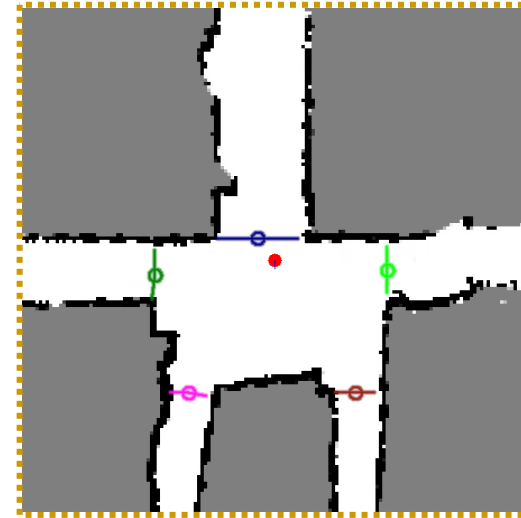
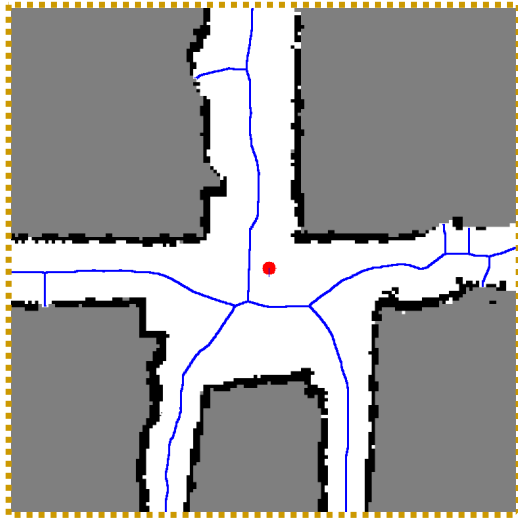


# Finding Gateways

- “A gateway occurs where there is at least a partial visual separation between two neighboring areas and the gateway itself is a visual opening to a previously obscured area.”
  - Chown et al., “PLAN” paper, 1995
- In previous work, we used gateways to determine the local topology at places.
- Multiple valid gateway definitions exist.
  - See paper for our formal definition.
- Gateways can also be determined at non-places.



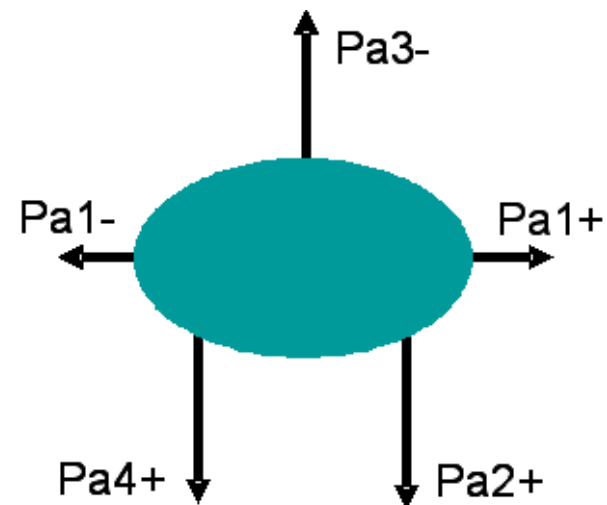
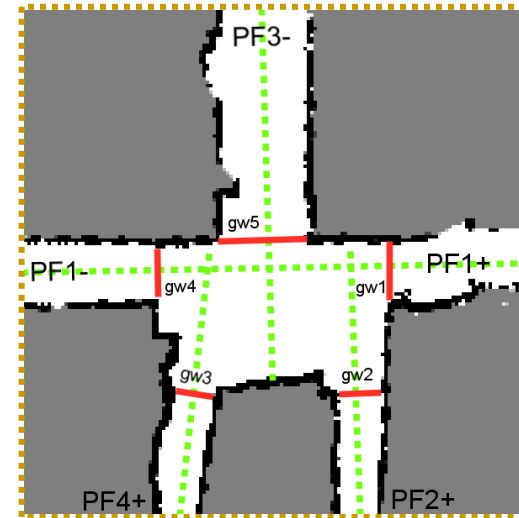
# Gateways in Complex Places



- Our gateway algorithm even works in complex intersections.

# Defining Paths

- Each gateway is associated with exactly 1 path.
- If exactly 2 gateways are unambiguously aligned, they belong to the same path.
  - Path passes through the place.
- Otherwise, that path terminates at the place.



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# Detecting Places

- The robot **IS NOT** at a place when:

- Exactly 1 path

**AND**

- Exactly 2 gateways

- Otherwise, the robot **IS** at a place.

- enclosed places

- Number of paths = 0

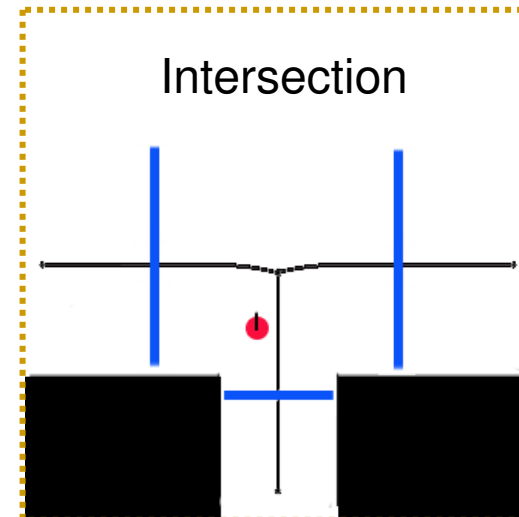
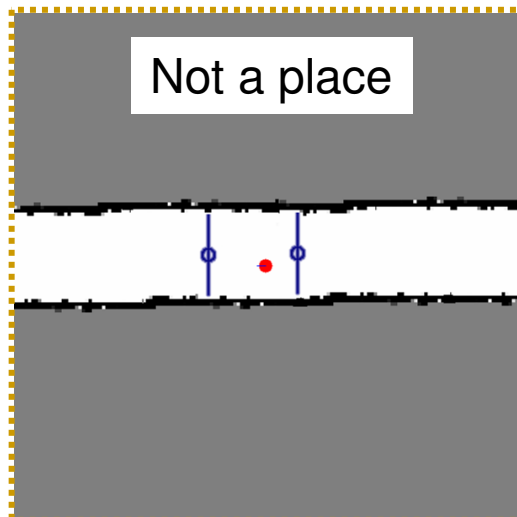
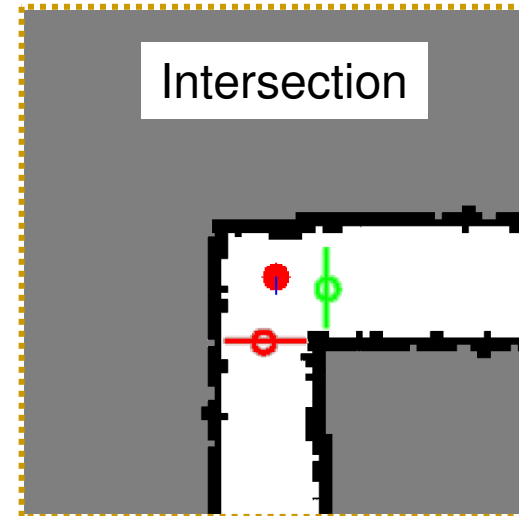
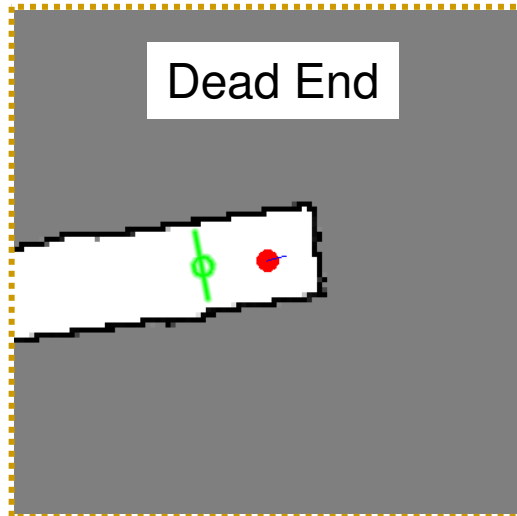
- dead-ends

- Number of gateways = 1

- intersections

- Number of paths > 1

# Place Detection Examples





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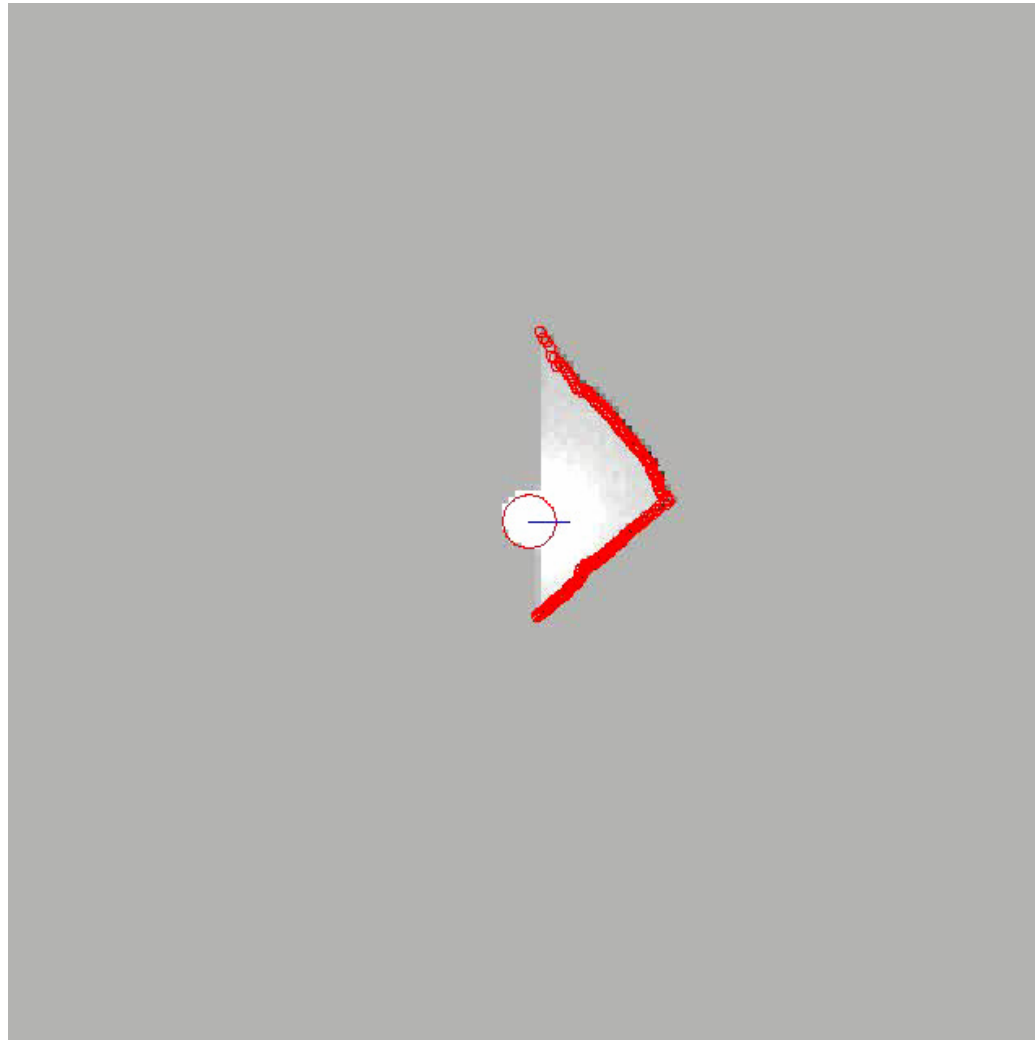
# Conclusion

- Our solution achieves reliable autonomous topological place detection (finding intersections, dead-ends, and open doors).
- Our method handles more environments.
  - Extended Voronoi graph handles both enclosed and non-enclosed environments.
- Previous false positive and false negative place detections are overcome.
  - Better pruning technique
  - Define places using gateways and paths not Voronoi junctions
  - Using our gateway algorithm, a robot can even detect complex intersections as single places with non-trivial local topologies.

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# Thank You

<http://www.cs.utexas.edu/~robot>



Video